Thesis Dissertation

IMPLEMENTATION AND EVALUATION OF ALGORITHMS ON CLASSICAL AND QUANTUM COMPUTERS

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Implementation and Evaluation of Algorithms on Classical and Quantum Computers

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Abstract

Quantum computing is an emerging technological breakthrough and it is making strong advancements towards commercialization. Its scope has shifted from the scientific community to a wider audience including major industry organizations such as Google, Amazon, IBM, organized communities interested in Quantum development, and finally independent developers. However, the quantum computing community is still in its early stages and there is a lot more improvement that could be made in the educational aspect concerning resources available and general approach.

This thesis aims to provide a streamlined workflow for Quantum Computing Development with the usage of the IBM Quantum Experience and Qiskit platform. The "Quantum Computing Learning Gate" was developed as a repository on GitHub consisting of four different levels of complexity. The first level is intended for experimenting with Quantum primers. The second level provides an interactive experience for understanding quantum computing behavior through the implementation of "The Exciting Game", a two-player game that is based on quantum mechanics concepts. The third and fourth levels complete an Implementation and Evaluation of the Deutsch-Jozsa and Bernstein-Vazirani algorithms on quantum and classical systems. In the evaluation of Deutsch-Jozsa, we have found that the quantum implementation shows an exponential speedup over the classical implementation with sufficient input size, however the accuracy of the current quantum systems is not sufficient to produce the correct outcome for the quantum solution since the correct result was found less than 1% of the time. In the evaluation for Bernstein-Vazirani, we have found that the current quantum systems are far from achieving a speedup over the classical implementation, but that the quantum accuracy of the algorithm is much higher than Deutsch-Jozsa, achieving accuracies close to 40% for the largest possible input on a quantum system.

Overall, we conclude that, while Quantum computing is still in its infancy commercially, there can be serious benefits to getting acquainted with the current stateof-the-art quantum frameworks, especially since quantum services are available through the cloud and can be a viable solution to real-world problems.

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Chapter 1

Introduction

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1.1 Motivation

The reasoning behind the existence of Computing is solving problems better than humans. Better meaning more efficiently, faster, and with fewer errors. The first "modern" computers were developed in the 1950s and we have made great strides in computation since then. However, as our calculative force increases, so does our curiosity and desire to conquer harder problems. Conventional machines, at their strongest form, can now output a staggering performance of upwards of **10¹⁷ Floating Point Operations**

per second [22],



Figure 1-1 Fugaku Supercomputer - by Fujitsu.

and the time we enter Exascale Computing is only some years away. We have created machines that can ridicule any given every day intensive calculation, yet for some problems, we can still make them calculate until the end of time and then some. A famous example is prime factorization which given a large enough number can prove surprisingly time-consuming. The best-known algorithm that solves the prime factorization problem

scales in sub-exponential time and as we know we can make some pretty big numbers considering the fact they are infinite. Now when we begin to look at some more specific problems it can get somewhat overwhelming. Examples of problems that can't be quite contained in a classical environment are quite more common than we would think, especially in the scientific community.

The simulation of a Quantum System for example or complex molecular structures found in Chemistry are two major examples that show the weakness of Classical Computers. In Newtonian mechanics, we can describe a system in a pretty discrete manner. However when we want to describe the dynamics of a Quantum System and we have to take into account the "magical" concepts of "Superposition" and "Entanglement" we are looking at a behaviour where two quantum objects can be perfectly correlated even though when looked individually can be described as random. An example given to often describe entanglement is that of a coin toss. Suppose we have two double-sided coins. When we measure the first coin it's up roughly half of the times and down the rest. However, the measurement of the second coin is always the opposite of that of the first coin. Already we are dealing with some non-classical properties. If we increase the number of "coins" or particles it can be shown that the number of all possible combinations of such a system can prove to be an intractable problem for a conventional machine [14].

These types of arguments were formally proposed to the Computer Science field by renowned physicist Richard Feynman, who urged the Computer Science community to take advantage of these concepts to try and move from theory to practice [7]. The initial motivation for the notion of Quantum Computing is these complex problems, which cannot yet be fully implemented on a practical scale since Quantum Computing is still in infancy and still heavily debated as to its contribution capabilities as stated here [10] and here [6]. However, the number of people who can experiment with these concepts in their simplest form and potentially contribute to the development is vastly bigger than what it used to be in the fifties.

With emerging technologies such as Cloud Computing Services, which give access to real quantum systems, communities of like-minded people who are curious about the elusive concepts behind what makes Quantum Computing work, we can argue that post-quantum cryptography and all the other applications are closer than we think. Although it seems to be that these applications are limited it is very likely that there will be some form of commercialized use in areas such as Cybersecurity, Weather Forecasting, Drug Development, and Artificial Intelligence. With behemoths such as IBM, Google, Microsoft, Amazon, JP Morgan Chase, Volkswagen Group already seriously investing in Quantum resources it does not seem like a bad idea to educate ourselves in this field.



Figure 1-2 IBM Q System One, The first Integrated Quantum Computing System for Commercial Use.

For this undergraduate thesis, our goal is to familiarize with the State of the Art Quantum technologies that are available to software developers and share our findings in the methodology of implementing some Quantum Algorithms while making practical use of the physical resources available through the cloud of IMB Quantum Experience. Along the way, we will share our understanding of some famous algorithms and concepts from a programmer's point of view, in our journey to prepare for the upcoming "Quantum Revolution."

1.2 **Thesis Overview**

Quantum Computing is an elusive concept, not only regarding Information Science but in Physics as well. This thesis is multi-purposed. Our objectives are as follows:

- Demystify some of the "simpler" concepts of Quantum Computing with pragmatic examples.
- Demonstrate a practical workflow that allows a somewhat adequately experienced individual in terms of pure Programming capability and Internet Technologies to start experimenting with Quantum Computing and avoid common pitfalls with the creation of the Quantum Computing Learning Gate which we will introduce in Chapter 4.
- Utilize our findings in order to have a better understanding of some of the more famous Quantum Computing Algorithms in order to implement them and evaluate them against classical implementations.
- Demonstrate the peculiar functionality of Quantum computers by making a minigame that is solely based on the properties of Quantum Mechanics.

Before proceeding with Qiskit experimentation we must state a few disclaimers as well. Below we provide what is **not** a goal of this thesis.

- Qiskit is a very elaborate framework that consists of hundreds of dedicated instructions, [32], and the purpose of this thesis is **not** to provide an introductory course to Qiskit usages and features. The reader is strongly advised to study the official Qiskit material which is provided here [34].
- This thesis is not claiming to **rigorously** prove physical proofs of why these Quantum Systems work or how they work. Whatever worked example we provide, is strictly based on our level of understanding and for the sake of explaining algorithm behaviours.

During our research, it was quickly made clear that quantum computing is not a final solution to all problems of Computer Science [10], [13], [6]. However, since so many behemoths of the industry have chosen to move quantum computing forward, we, as computer scientists need to raise a number of research questions.

The first research question that we raise, concerns the crowd of quantum computing. Is quantum computing for everyone? Is it a subject that should be compulsory in every university's curriculum?

The second research question is about which are the current State-of-the-Art resources available for quantum computing and which one we should choose depending on our level of expertise or needs. During our research, we had to choose a medium in which we could conduct our experiments and acquire valuable information in order to educate ourselves on a completely new subject.

The third and final research question is whether or not a career specialized in quantum computing is a viable option in the next 5 to 10 years. This is of course a question that is dependent on many different factors. Either that be major technological breakthroughs or agenda shifts of major corporations it is still something that we need to ask ourselves.

1.3 Thesis Contribution

With the completion of this thesis we achieve the following milestones:

The first milestone is the completion of a functional education portal which we call "Quantum Computing Learning Gate" or "QCLG". Through the development of this repository in GitHub, we can offer a no-nonsense approach to quantum computing education through a variety of different experiments, algorithm analysis, evaluation, and interactive tools.

The second milestone is the creation of "The Exciting Game", an interactive twoplayer game, incorporated in the QCLG platform, which can help individuals familiarize themselves with basic quantum concepts through a gaming experience.

The third milestone is a comparison study of algorithms on quantum and classical computers. From this, useful data was extracted concerning two major algorithms in the field of quantum computation by simply building upon the basic blocks available through the Qiskit platform.

The fourth milestone is the expansibility of our platform. Not only can one learn from QCLG, but they can also contribute to any of the different levels of development available in QCLG.

The fifth milestone is that our platform can be a central point for further quantum computing exploration since inside QCLG, we list some different sources that we have found extremely useful throughout our own research.

1.4 **Thesis Outline**

In the first chapter, we roughly presented the incentive of this thesis and talked about the current picture of Quantum Computing. We also raised some research questions that will be answered by the end of this thesis. In the second chapter, we will talk about the various frameworks available for quantum computing and showcase some interesting applications that are closely related to our work. In addition, we will discuss the recommended background knowledge required. In the third chapter, we discuss the basic components that help us construct more complex algorithms. In the fourth chapter, we provide a high-level description of the structure of QCLG. The next chapters will directly expand on the structure given in this chapter. In the fifth chapter, we dive into the first level of QCLG and execute and analyse some simple experiments that underline the major quantum properties we will use in the rest of this thesis. In the sixth chapter, we take an extensive look into the second level of QCLG which is about "The Exciting Game" and how it can help boost the usefulness of the whole project. In the seventh chapter, we showcase level 3 of QCLG by implementing in detail both classical and quantum approaches to two famous algorithms, Deutsch-Jozsa and Bernstein-Vazirani. In the eighth chapter, we showcase the final level of our platform which is the evaluation of algorithms. Here we implement a series of automated tests in order to extract useful data about the two algorithms. In the ninth chapter, we present our technical conclusions concerning the development of programs with Quantum modules and discuss limitations and future work.

Chapter 2

Related Work and Background

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2.1 Related Work

Since this thesis is multi-purposed, there is a number of different works that are close to our efforts.

2.1.1 IBM Quantum Experience and Qiskit

The first and major example is the Qiskit framework and textbook, which presents an abundance of information regarding basic quantum concepts, along with worked examples and dives into deeper and more complex problems. The largest portion of this thesis was inspired by the Qiskit textbook. Qiskit is an open-source framework that started in March of 2017. Developed by IBM Research, it now has a strong following with the Qiskit community and has a public repository on GitHub. Qiskit is mainly developed in Python and is cross-platform compatible. The Qiskit textbook is a university-level supplement to quantum computing courses. It offers the ability to work with machinelevel code with OpenQASM, [4], as well as high-level abstractions using Python. All of Qiskit is encapsulated in the IBM Q platform, [25]. IBM offers free experimentation on real quantum systems, as well as simulators that work under the Qiskit framework. It also offers interactive tools with no need for coding and a plethora of information through a very well-built documentation library, [26].

2.1.2 Amazon Braket

Amazon is a giant in the industry and it is not a surprise that it now offers an elaborate cloud quantum computing service for researchers and developers [16]. However, Amazon Web Services in general, require a certain level of specialized expertise, which was impossible to achieve in an academic semester in order to accomplish our goals for this thesis. It is worth noting that Amazon uses D-Wave systems.

2.1.3 **D-Wave - Leap**

D-Wave was founded in 1999 and is a commercial supplier of quantum computers. D-Wave systems are used by the leading technology organisations including Google, NASA, Volkswagen. The research that stems from D-Wave is very rich with over 100 peer-reviewed papers in scientific journals. In 2018 D-Wave launched the Leap, their own cloud quantum computing service which contains open-source educational material meant for developers, researchers, and businesses. Leap² is directed to people

looking for hybrid solutions that tackle problems with the best mix of quantum and classical resources. It was made for the development of hybrid quantum applications that are not only directed in scientific experimentation but real-world problems, [27],[21].

2.1.4 Quantum Tetris

In the course of work, we managed to implement a quantum game in its infancy. A motivation for this was Quantum Tetris [35]. It is about the famous game but with a quantum spin. It is implemented using IBM resources and Qiskit. While it is quite a different game than what we will present it opened our eyes concerning the simple applications that quantum computing might have.

2.1.5 **Procedural Game Generation**

As discussed here, [15], there are major efforts to start producing procedural generation for games and more using quantum computing. Concepts such as blur seem to have a better place in the hands of the stronger quantum computers that will come.

2.1.6 Volkswagen Quantum Routing

Volkswagen was the first company to demonstrate successfully a real-world application for quantum computers. Through demonstration on "quantum buses", Volkswagen presented a live use of a system designed to calculate the fastest routes in almost-real time using D-Wave's systems [38].

2.1.7 Next-Gen Batteries Using Quantum Computers

In a research paper conducted by joint efforts of IBM and Daimler AG [11], the parent company of Mercedes-Benz, simulations were conducted, of molecules that are formed in lithium-sulfur batteries while operations. This was demonstrated using IBM Q premium-access systems. The goal of this research is to pave the way for the design of the next generation of batteries installed in electric vehicles. We can see that more and more industries have a direct benefit from the development of quantum computing.

2.2 Background

2.2.1 Methodology

For this thesis, we have worked with IBM'S Quantum Experience resources. The operating systems we used are Ubuntu 18.04 and Windows 10, the programming language is Python and the main external library used to implement quantum programs is Qiskit. We have deployed a repository in GitHub and keep all our code. IBM Q offers extensive tutorials in order to set up a local environment capable of accessing IBM'S systems. It is important to note that the default medium for conducting quantum experiments is either's through Quantum Experience's interactive tools such as the Quantum Lab, where the user can design from scratch and execute a Quantum Circuit, or through Jupyter Notebooks where you can write code in normal Python with the necessary instructions from Qiskit in order to implement the Quantum Algorithm. For this thesis, we chose to work with PyCharm since it's a very popular Integrated Development Environment and good practice for testing out IBM's API outside of the recommended frameworks.



Figure 2-1 PyCharm.

Q

Figure 2-2 IBM Quantum.

 \mathbf{O}

Figure 2-3 GitHub.

2.2.2 PyCharm

PyCharm is an Integrated Development Environment created by JetBrains. Aside from being a great Python IDE it also possesses multiple versioning control features with built-in GitHub support as well as Conda environments More information can be found here [30].

2.2.3 Conda environments

In order to maintain universality across Operating Systems, we opted to work with Conda environments. They provide the ability to work with a Python interpreter which can be used in different Operating Systems, [18]. Our Conda virtual environment contains:

- Python 3.7
- Qiskit
- NumPy (version: 1.19.3)
- Matplotlib

This environment can easily be created using the online instructions provided in the conda documentation and then proceed to install the aforementioned dependencies using pip install. Any clarifications for Ubuntu commands can be found here [37].

2.2.4 **Python 3.7**

Python is an open-source, interpreted-style, programming language that is widely used both in the industry and in academia, because of its simplicity and open-source nature. It is also the underlying language that Qiskit uses. For this thesis, in order to create a stable environment where multiple dependencies coexist, we used Python 3.7, [31].

2.2.5 NumPy

NumPy is one of the dependencies required by Qiskit and one that we had to install manually to ensure a stable working environment. Qiskit being dependant on NumPy makes a lot of sense, since as we will see later, Qubits and System States, in general, are described by vectors and matrices, [29].

2.2.6 **GitHub**

GitHub is a famous code hosting platform for version control using Git. We used GitHub in order to ensure a clean workflow and easily transfer our work locally to another device, [23].

2.2.7 Matplotlib

Matplotlib is a very strong library for creating visualizations in Python. This was especially useful since Qiskit supports a number of different visualizations and proved vital in order to debug Quantum Circuit construction and help us understand the algorithms' structure in a more interactive way, [28].

2.2.8 **IBM Quantum Platform**

The IBM Quantum Platform has a lot of useful tools that can help with quantum development, without the use of code. An excellent tool is the circuit composer, [24], which is a very nice way of experimenting with simple experiments, as well as complicated algorithms. We present a more detailed overview of the circuit composer in 3.3.2. The interface of the IBM Quantum Platform is very rich. Besides the circuit composer, the user can view the available quantum systems, access relevant documentation, and develop quantum programs using different tools.

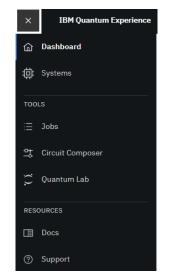


Figure 2-4 IBM Quantum Experience Interfaces.

Welcome		Recent circuits	(35)			Your systems (9)
Sotiris Loizi	dis	Name		Last update	d	Maintenance alert:
Your providers Personal profile		Untitled circuit	Untitled dirout I days ago			In preparation for launching 3 new 5-qubit systems on Feb 10th, ibmq_ourense, ibmq_vigo and ibmq_valencia
		Untitled circuit				will be retired Jan 20th. As a result, users may experience longer queue times during the transition.
See more		Untitled circuit				ibmq_armonk will be offline around Jan 4th-26th. We apologize for any inconvenience this causes and hope
Experime	nte	Untitled circuit		a month ago		you will like the new systems.
-				a month ago		Online
awaiting	executio			1000100		ibmq_santiago (5 qubits, QV32)
1		View All				Total pending jobs: 16 jobs
Pending results (0)			Ava	ilable Syster	ns, depending	Orline
You have no circuit runs in t	he queue			ccount right		ibmq_athens (s qubits, QV32)
Tou have no circuit runs in t					5.	Total pending jobs: 23 jobs
	La	atest completed	1			() Online
Latest results	- ex	periments.				ibmq_valencia (5 qubits, QV16)
Status 💠	Run date 👻	Name or Id	Tags	Provider	Service ¢	Total pending jobs: 386 jobs
COMPLETED	3 days ago	5ffa3f8950c62cfffe9c8900		ibm-q/open/main	Backend: ibmq_valencia	
COMPLETED	3 days ago	6ffa3ecd50c62c303a9c88f3		ibm-q/open/main	Backend: ibmq_valencia	ibmq_vigo (5 qubits, QV16)
COMPLETED	3 days ago	5ffa3d4a935c00c1e5212029		ibm-q/open/main	Backend: ibmq_valencia	Total pending jobs: 700 jobs
COMPLETED	3 days ago	5ffa371d6d095068a16caad0		ibm-q/open/main	Backend: ibmq_valencia	() Online
COMPLETED	3 days ago	6ffa32d63d08740d4ba20917		ibm-q/open/main	Backend: ibmq_athens	ibmq_16_melbourne (15 qubits, QV8)
COMPLETED	3 days ago	5ffa3064234c5d06ad9eafdf		ibm-q/open/main	Backend: ibmq_athens	Total pending jobs: 711 jobs
COMPLETED	14 days ago	eval_13_	ibmq_jobset_f820	ibm-g/open/main	Backend: ibmq_16_melbourne	Online
COMPLETED	14 days ago	eval_12_	ibmq_jobset_f820	ibm-q/open/main	Backend: ibmq_16_melbourne	ibmq_ourense (5 qubits, QV8)
COMPLETED	14 days ago	eval_11_	ibmq_jobset_f820	ibm-q/open/main	Backend: ibmq_16_melbourne	Total pending jobs: 794 jobs
COMPLETED	14 days ago	eval_10_	ibmq_jobset_f820	ibm-q/open/main	Backend: ibmq_16_melbourne	🚯 Online
View All						ibmq_5_yorktown (s qubits, qv8)
						Total pending jobs: 543 jobs
						Online
						ibmq_qasm_simulator (Up to 32 qubits)
						Total pending jobs: 56 jobs
						8 Maintenance
						ibmq_armonk (1 qubit)
						Total pending jobs: 1 job
						vites permits jour + jou
						V* Polse

Figure 2-5 IBM Quantum Experience Dashboard

In Figure 2-5 we can see useful information regarding the available systems, as well as a history of all experiments. Clicking on one of the systems, the user can see information regarding a system's hardware.

In the Quantum Lab, shown in Figure 2-4, the user can create Jupyter Notebooks, which is the default way for Quantum development using Qiskit.

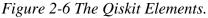
In the Jobs label we can see the progress of completed experiments or experiments waiting in the systems queues.

We have implemented necessary methods to bypass long queues by finding the least busy backend capable of executing our experiment. These methods can be found in Appendix E.

2.2.9 Qiskit SDK

Qiskit is an open-source Software Development Kit that is used for development using Quantum Computers. It can be used to develop low-level Quantum Instructions by tinkering with pulses or to develop higher-level applications using Quantum Circuits and complete application modules.





From the picture above we can see that Qiskit is divided into four elements. Each element has its purpose. Terra is the base for all other elements. It contains the modules necessary to develop on the lowest level, modules necessary for parameter tuning in different hardware as well as modules that are responsible to handle the completed circuits as executable jobs to the designated systems. Aer is responsible helps as a benchmark of the Quantum Computing efforts by providing a realistic simulation of Quantum Systems using classical computers. Ignis is an effort to combat error in quantum systems by providing different ways of examining quantum circuits, gates, or code. Aqua is focused on Quantum Computing Applications. It requires expertise in Quantum Computing and was created for domain-specific experts. More information can be found here [1], here [32], and here [34].

In our thesis, we mainly focus on the modules provided by Terra and Aer. We list the most useful to our thesis:

Terra:

- qiskit.circuit, A model where operations or qubits are performed using quantum gates.
- qiskit.providers, Useful for interacting with circuits that are currently running.
 - Provider: Provides access to different backends, according to provider rights.
 - Backend: abstractly represents the real system or simulator and is responsible for executing the quantum circuits and returning the results.
 - Job: A job is essentially the key for a certain experiment. It is useful for keeping track of the progress of the experiment, queued, running, etc, and providing the ability to control it.
 - Result: An object which keeps the quantum data of a remote backend from

 a completed experiment. It can be instantiated using result = job.result(). There are multiple ways to manipulate this data but the
 most common is through result.get_counts(circuit).
- qiskit.visualization, Contains multiple tools for visualizing quantum components.

Aer:

QasmSimulator: Powerful Simulator that allows for the execution of experiments in ideal quantum computers or even environments close to real systems. For this thesis, we use it for its ability to host ideal circumstances.

2.2.10 Linear Algebra

Quantum computing requires a certain level of experience in linear algebra. From complex numbers to vectors to matrices, to multiplication and addition of different elements, there is a lot to learn if one wants to understand better how quantum systems behave. It is probably the most useful tool in order to have a real understanding of the qubit states and how they change. More information on linear algebra for quantum computers can be found here [2] on page 17.

2.2.11 Classical Gates

Quantum computing uses the quantum gates as abstractions to the actions performed on qubits. Therefore, although classical and quantum computation have major differences, we can see that there is some common ground. Many of the gates used in Qiskit are similar to traditional classical gates.

One example is the classical NOT gate which flips the state of a bit. In quantum computing, there is the X gate which flips the state of a qubit.

Another example is the XOR gate which acts based on two bits. The corresponding gate of quantum computing is the CX gate or C-NOT gate.

An individual would certainly find it easier to experiment with quantum gates if they revisit the classical gates [9].

2.2.12 Basic Quantum Mechanics

The hardware of quantum computing is nothing like a classical computer. It is advised that the aspiring quantum programmer acquires basic knowledge of the materials needed in order for a quantum system to exist. More can be found in [2], chapter 1.

Chapter 3

Quantum Computing Primers

3.1 Ke	y Concepts	17
3.1.1	Superposition	
3.1.2	Entanglement	
3.1.3 Quantum Interference – Phase Kickback		
3.2 Qu	bit	
3.2.1	Bloch Sphere	
3.3 Qu	antum Circuits	
3.3.1	Quantum Gates	
3.3.2	Circuit Composer	
3.3.3	X gate	
3.3.4	CX "CNOT" Gate	
3.3.5	Hadamard Gate	

3.1 Key Concepts

Quantum Computing has some key concept differences from classical Computation. More information concerning these concepts can be found here: [2].

- Superposition
- Entanglement
- Quantum Interference

3.1.1 Superposition

Superposition is something that is not a property of Newtonian mechanics. Superposition is the ability of a Quantum System to be in a combination of multiple states at once. Suppose we have a Quantum Computer with n qubits, there can be a superposition of all possible 2^n states at once. This property is useful in many Quantum Algorithms.

3.1.2 Entanglement

Entanglement is perhaps one of the most curious concepts in the world of Quantum Mechanics. It is closely related to superpositions and it tells us that two or more particles are in such an assortment that their quantum state cannot be described independently of the others in that same "correlated" group. Abstractly speaking, the behaviour, of one particle will directly affect or be affected by the behaviour of some other particles that are somehow related to that particle.

3.1.3 Quantum Interference – Phase Kickback

A fundamental idea in Quantum Computing is being able to control the probability of measuring certain states. It is a by-product of the superposition and it allows us to tinker with certain superpositions in order to steer the system in a deterministic outcome which is useful to us.

3.2 **Qubit**

The bit is the simplest single point of information a Classical Computer can describe. If we strip away all ingenious hardware and software designs of the modern machines, a bit is what we are left with. It is the "atom" in the universe of traditional Computation. A bit holds a discrete value of either one or zero. The most used processors today are children of the 64-bit architecture. Explaining 64-bit architecture very roughly we, can say that at least theoretically an architecture of this magnitude that a 64-bit processor can address 2⁶⁴ bytes or 16 exabytes of byte-addressable memory, a memory that can be referenced for each byte with a different address. In Quantum Computing, our "atom" of computation can quite literally be an atom.

While a bit can be described with the functionality of the flip-flop, for example, the qubit is a bit more complicated. There is a number of ways that a qubit can be implemented but let's just imagine an electron. An electron has the property of spin, which we will not delve into, but it can translate to two separate states, spin up and spit down, thus creating a two-state quantum-mechanical system [2].

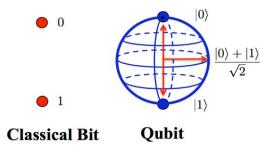


Figure 3-1 Quantum Optics and Quantum Many-body Systems Andrew Daley's Research Group at the University of Strathclyde [36].

What differentiates the qubit from the bit is the additional physical properties acquired by quantum mechanics. Quantum mechanics allow the qubit to exist in a state of both 1 and 0 simultaneously. This property of the qubit is fundamental in quantum computing.

3.2.1 Bloch Sphere

Named after physicist Felix Bloch, the Bloch sphere gives a geometrical representation of the qubit. From the figure below we can observe three axes. X, Y, and Z. The north and south poles of the system as chosen as the ground and excited states of the qubit $|0\rangle$ and $|1\rangle$ and correspond to the spin-down and spin-up of the electron. Points on the surface of the sphere are defined as pure states, states that when are measured can be expected to yield a certain value, and points within the surface of the sphere are defined as mixed states or states that are a statistical ensemble of pure states, which essentially means that they are in a state that is a "superposition" of pure states and we cannot be certain on which of the pure states we will end up measuring. It is worth noting that the "amount" of a pure state in a mixed state configuration is determined by a complex number.

These states are described using Dirac notation. The general case for the two basis states of $|0\rangle$ and $|1\rangle$ is:

 $c_0|0\rangle + c_1|1\rangle$ where c_0 , c_1 are complex numbers.

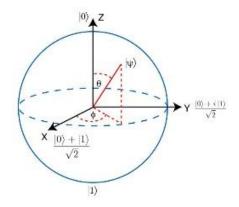


Figure 3-2 Bloch Sphere from the Qiskit Textbook [17].

3.3 Quantum Circuits

Quantum circuits are a model in which a computation is performed with a sequence of quantum gates, which are reversible transformations on a quantum mechanical comparable of an n-bit register. This analogous structure is referred to as an n-qubit register. Further operations can be performed on a quantum circuit like logical operations, barriers in order to control the circuit compilation, and of course measurements in order to get results.

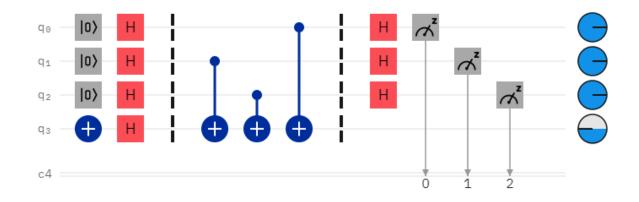


Figure 3-3 A Circuit Example, Containing Multiple Quantum Gates, and Operations.

3.3.1 Quantum Gates

The reader should be familiar with the classical gates of computation such as the AND, OR, XOR, NAND, and so on. The purpose of quantum gates is to set the qubit into a certain physical state so that we can dictate with certainty or to some degree the outcome of the measurement of the qubit or in order to further develop the state of the quantum system as a whole. Quantum gates can be unary operators, acting on a single qubit, or even operators that include multiple qubits in order to be implemented. What quantum gates essentially do, is rotate the qubit to a state. All quantum gates can be constructed by programming them to "rotate" the qubit on any combination of the three axes shown in the Bloch sphere in order to reach a distinct state. There is an infinite number of different physical states that we can shape but below we show a few significant gates that help us on the rest of this thesis. Quantum gates are reversible. Thus, they can be represented as matrices, manipulating the state vector of the Bloch Sphere. Each time a gate operation is performed on a qubit, we can calculate the new state of the qubit by matrix and vector multiplication to find the new state.

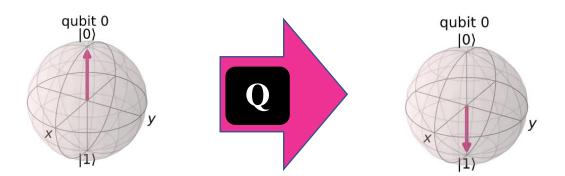


Figure 3-4 A Qubit State Transformation Through a Quantum Gate Application.

One important note is that we can perform operations on different bases, the three mainstream bases being X, Y, and Z.

Looking at

Figure 3-4 we can see that we could have an infinite number of bases by simply using two orthogonal vectors each time. For all experiments shown in this thesis, we will use the Z-basis, which is the most popular.

The two basis states, $|0\rangle$ and $|1\rangle$ are represented in a vector like so:

$$|0\rangle = \begin{bmatrix} 1\\ 0 \end{bmatrix} \qquad |1\rangle = \begin{bmatrix} 0\\ 1 \end{bmatrix}$$

3.3.2 Circuit Composer

A very helpful way of understanding and experimenting with quantum gates is the Circuit Composer tool, provided by IBM Q, where we can interactively place quantum gates on a "quantum circuit" and perform all kinds of operations including, gate operations, classical logical operations, measurements and more.

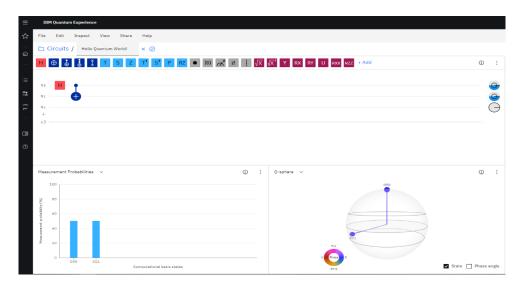


Figure 3-5 A More Comprehensive View on the Circuit Composer.

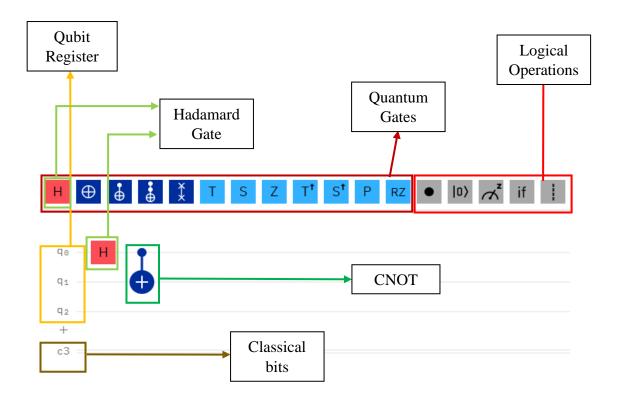
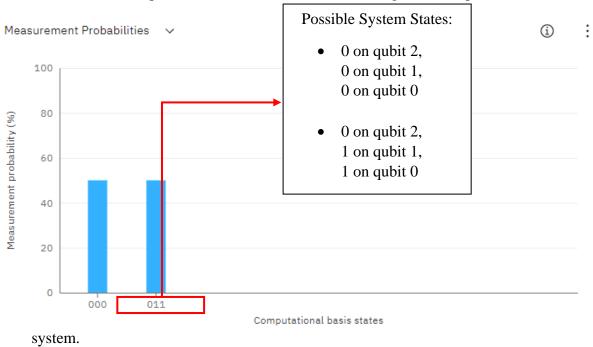


Figure 3-6 A Detailed View of Circuit Construction.



The measurement probabilities section showcase a histogram of all possible states of the

Figure 3-7 Measurement Probabilities.



Figure 3-8 Q-Sphere Representation.

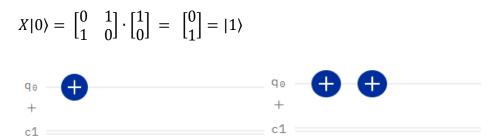
The Q- sphere representation is a very nice way of visualizing the global scale of the system. There are a lot of compact details about the system in this single frame. It is different than a Bloch sphere visualisation since it shows the possible states of the system of qubits as a whole. We can also examine the likelihood of each system state and the angle that the phase on the measuring base.

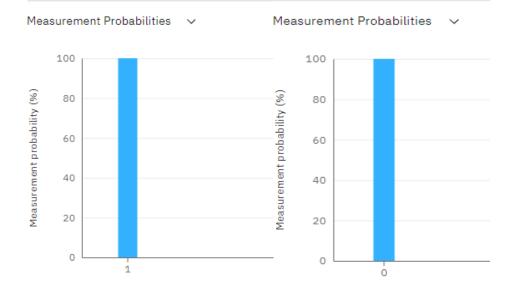
3.3.3 X gate

The NOT gate of Quantum Computing. The X gate acts on a single qubit and switches from state $|0\rangle$ to $|1\rangle$ and vice versa.

The X matrix: $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

Applying an X gate to a qubit that is in state $|0\rangle$ will then have the following outcome:



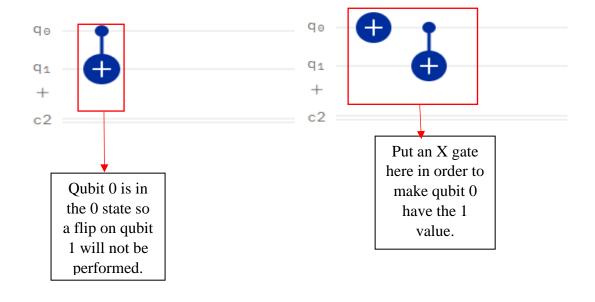




On the right side of Figure 3-9, we can see the reversibility of gates since by applying two X gates one after the other we return to the 0 gate.

3.3.4 CX "CNOT" Gate

The controlled-NOT gate is a binary qubit gate since it needs two qubits to operate on. One target qubit and one control qubit. Whenever the control qubit is in the $|1\rangle$ state and X operation is acted on the target qubit. Later we will see how this gate helps us achieve entanglement.



Measurement Probabilities 🗸 Measurement Probabilities

Measurement Probabilities 🗸 🗸

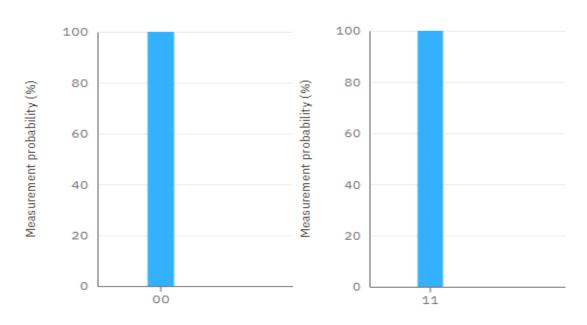


Figure 3-10 CX or C-NOT Gate Behaviour.

3.3.5 Hadamard Gate

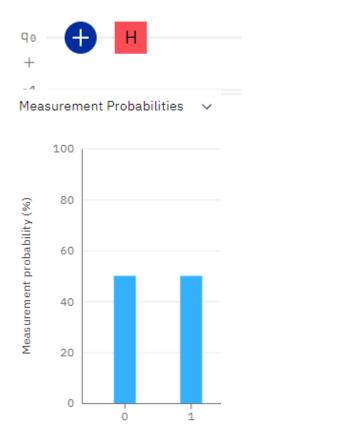
The Hadamard gate is a very special operator. It is useful when we want to achieve superpositions entanglement or interference. It will be the cornerstone of all major experiments in this thesis.

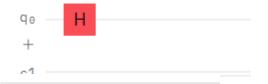
The H Matrix:
$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

Applying an H gate to a qubit that is in state $|0\rangle$ will then have the following outcome:

$$H|0\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = |+\rangle$$
$$H|1\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \cdot \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} = |-\rangle$$

Both $|+\rangle$ and $|-\rangle$ are states which are essentially in a superposition of exactly equal probability of yielding 0 or 1.





Measurement Probabilities 🗸 🗸

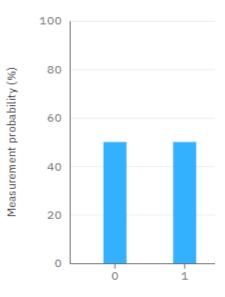


Figure 3-11 Hadamard Gate Behaviours.

Chapter 4

Quantum Computing Learning Gate

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4.1 Introduction

Once an individual is familiarized with the quantum primers, it is time to move forward. We have developed a repository in GitHub in our efforts to create a standalone platform capable of hosting experiments of different levels. All necessary instructions to deploy this project can be found on Appendix F.

4.2 Structure Breakdown

After many different designs, we have finalized the structure for this platform with simplicity in mind. The idea of QCLG is to allow individuals of different competence levels to experiment with Quantum Computing. It is also possible to contribute to different levels according to the structure rules. We have implemented four different levels of experimentation, each with a different goal. All levels take advantage of Qiskit resources but to different degrees.

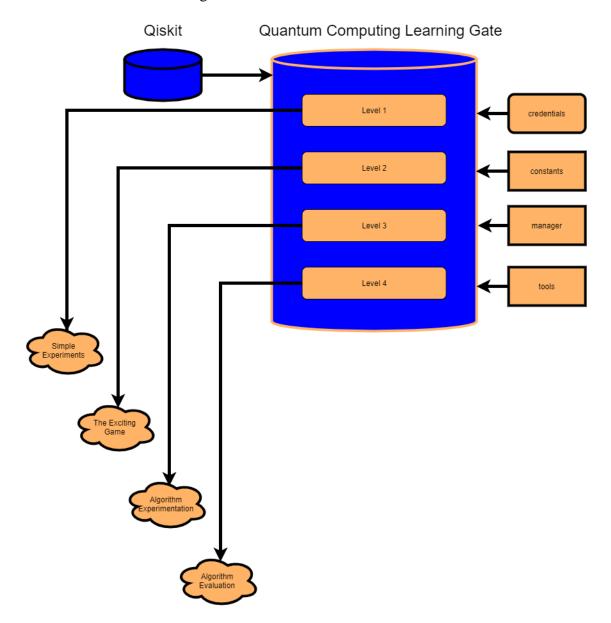


Figure 4-1 The QCLG Platform.

The first level aims to cement the basic knowledge required in order to proceed with more advanced concepts. Based on our learning experience, we believe that an intermediate step was required before tackling quantum algorithms. Although very interesting, quantum algorithms can prove to be quite uninviting to beginners. We decided to construct a level where the user can fall back on and sharpen their understanding by playing quantum games, hence the creation of Level 2. This way they can try again experimenting on the next level with more confidence. Level 3 analyses the more mainstream quantum algorithms by attacking the problem from both sides of the spectrum. Here we provide implementations for both classical and quantum approaches to solving certain problems. The final level is for evaluating algorithms. Once the user possesses a deeper understanding of an algorithm, they can stress test it and extract useful results performance and accuracy-wise.

4.2.1 Level 1 – Experimenting with Primers

Level 1 is where experimentations concerning the basic concepts of Quantum Computing take place. It is the starting place after getting familiarized with the Quantum Primers we discussed in Chapter 3. Here, the user can explore and add experimentations of very basic quantum circuits in order to test their understanding before moving on to the next levels. The currently available experiments in Level 1 are about Superposition, Entanglement, and Phase Kickback and we showcase them in more detail in Chapter 5.

4.2.2 Level 2 – The Exciting Game

Level 2 is an intermediate layer that aims to motivate the user to learn about Quantum Computing through more interactive mediums. Once the user has become more proficient with the basic quantum concepts, they can test their knowledge with a quantum game for example. The currently available game is "The Exciting Game" which we cover extensively in Chapter 6.

4.2.3 Level 3 – Implementation of Algorithms

Level 3 has two objectives. The first objective is to implement quantum algorithms and the second objective is to implement the classical solution to that algorithm in order to allow the user to achieve a deeper understanding of that algorithm. Level 3 has many choices regarding the execution of experiments. They can choose to solve the problem classically by executing the solution onto their local device and observe

the results, execute the quantum solution of the problem in the IBM backends or solve the problem with both approaches and observe the differences. Below we list the three directories included in Level 3.

4.2.3.1 quantum_algorithms

In this directory, we store the implementations of major quantum Algorithms, as well as a controller class that allows individual calling of each of these algorithms. Up to the current point of development of this thesis, we have implementations for the Deutsch-Jozsa, which can be found in 7.3.3, and Bernstein-Vazirani, which can be found in 7.6.3, offering execution for each algorithm on a simulator, on a real quantum system, on the local machine, executed classically, on both a real quantum system and the local machine for comparison.

4.2.3.2 oracles

In the oracles directory, we store all the oracle functions necessary for implementing the quantum algorithms. Currently, there are implementations for one possible oracle function of Deutsch-Jozsa which can be found in 7.3.3, and the implementation of the oracle function for Bernstein-Vazirani which can be found in 7.6.3.

4.2.3.3 classical

The classical directory contains all required classes for implementing the classical solutions of the Quantum Algorithms. This includes the logic for the classical algorithms, as well as some additional implementations for input generation. These can be found in 7.2.2 and 7.5.1.

4.2.4 Level 4 – Evaluation of Algorithms

Once the user feels comfortable with all the previous levels, they can proceed to Level 4, which is the automated evaluation of an Algorithm based on its classical and quantum solution. Here, besides an adequate understanding of Quantum Computation concepts, the user must also practice with the Qiskit API in order to extract the necessary data to extract useful results. More information on the evaluation can be found in Chapter 8.

4.2.5 Independent entities

In Figure 4-1 we can see there are four separate entities outside of the QCLG Platform. credentials, constants, manager, and tools.

4.2.5.1 constants

The constants entity is a python file containing a variety of different text messages as well as stored acceptable inputs for many different parts of the project. Its main purpose is to increase readability on the rest of the project.

4.2.5.2 manager

The manager entity is a python file that acts as the main driver of the whole project. From this point, we can branch to every different level available in this project.

4.2.5.3 credentials

This is where the user can store their IBM Q credentials to access the remote resources of IBM. In order to access the IBM resources an account in IBM Q is required. After creating an account, the user can generate a token that allows them to access IBM resources remotely.

4.2.5.4 tools

The tools entity is a python file and is the backbone of the whole project. Here we implement an amalgamation of methods, consisting of remote calls to IBM resources, calls to different parts of the project with customizations, evaluation methods, and more.

The project Tree

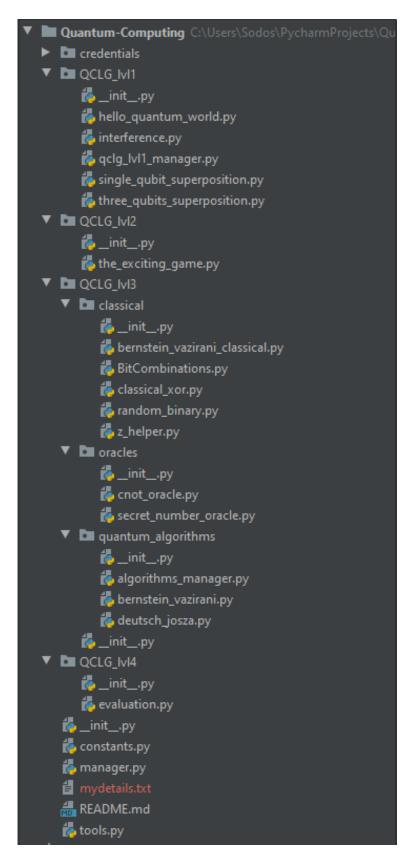


Figure 4-2 The Project Tree.

4.3 **GitHub Repository**

This project is hosted by the Data Management Systems Laboratory (DMSL) -Department of Computer Science, University of Cyprus, more information about DMSL can be found here [20].

ÿ;	Data Management Systems Laboratory (DMSL), CS, UCY
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Q Find a repos	itory Type: All • Language: All •
quantum Quantum Progra ● Python 栜 GF	mming 1-3.0 양 1 ☆ 1 ① 0 \$\$ 0 Updated 3 days ago

Figure 4-3 DMSL GitHub Page

The Quantum Computing Learning Gate is an ongoing project, and its current state of development, as well as previous or future states of the project, can be found here [19], under the "dmsl/quantum" label. This project is under the "GNU General Public License v3.0".

We have also supplemented Quantum Computing Learning Gate with a separate README file, containing the essential information about the project, and the deployment guide, shown in Appendix F.



We have decided to provide some additional useful resources that helped us during our research, since it can be quite difficult to find a good starting point.

Chapter 5

QCLG Level 1 – Implementing Primers

5.1	Introduction	
5.2	Achieving Superposition	
5.2.1	I Implementation	
5.3	The Bell State	
5.3.1	I Implementation	
5.3.2	2 Results	
5.4	Phase Kickback	
5.4.1	I Implementation	
5.4	4.1.1 Results	

5.1 Introduction

At this point, we start implementing a series of simpler experiments to practice on both the primers we discussed in Chapter 3 and methods imported from the Qiskit SDK as we presented briefly in 2.2.9. Some pointers to help out are to always keep a window open with the Qiskit documentation which can be found here [32] and the most important to press the control key + left click to be transferred to the class implementing each method in order to clear any doubt on what the current method achieves. We have supplemented our code with comments where it was deemed necessary, but the user should spend as much time as possible understanding both the concepts presented and the way they were implemented before moving on to the next level of QCLG.

5.2 Achieving Superposition

In this experiment, we will set up a circuit such that all qubits are transcended into a superposition using the Hadamard gate which will configure each qubit into a state with an equal probability of yielding zero or one and then measure the state of the system. We will execute this experiment 1024 times and observe the measurements. For this demonstration, we will use four qubits. The possible states with three are equal to $2^3 =$ 8. Since we use the Hadamard gate to set up our system we expect that each time we make a measurement we have an equal probability of measuring one of the 8 possible states:

$|000\rangle$, $|001\rangle$, $|010\rangle$, $|100\rangle$, $|011\rangle$, $|101\rangle$, $|110\rangle$, $|111\rangle$.

In an errorless quantum computing system, we expect to **approximately** measure each of the above states $\frac{1024}{8} = 128$ times. Of course, even in an errorless system, this should not be the case since what a Hadamard gate guarantees, is an equal probability of measuring either one or zero. Hence, we just expect to see all states measured to a number very close to 128. We can simulate the behaviour of an errorless system using the simulator provided by IBM Q. After assembling a three-qubit circuit with a Hadamard gate attached to each qubit and immediately measuring each qubit we get the two following results. The first result is achieved by executing the circuit on the simulator and the second by executing on a real device.

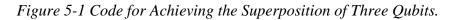
Simulator execution: Total count all possible states are: {'000': 134, '001': 129, '010': 102, '011': 127, '100': 119, '101': 129, '110': 132, '111': 152}

Real System execution: Total count for all possible states are: {'000': 185, '001': 167, '010': 114, '011': 92, '100': 152, '101': 111, '110': 87, '111': 92}

We can see that indeed we measure roughly as we expected. What we can also see is that while a superposition seems very powerful, it doesn't accomplish much if we are eager to measure qubits. Measurements in quantum computing are irreversible on the state of a system and we should not be measuring whenever.

5.2.1 Implementation

We implement the quantum circuit with the help of Qiskit. Qiskit offers a plethora of different classes, each for a different purpose. Some things to notice in the following code is that we import different classes from Qiskit



The major processes accomplished in this code is the creation of a circuit with 3 qubits and 3 bits with the call to the QuantumCicuit class. Then we execute the circuit both on the simulator and the ibmq_valencia backend using the necessary Qiskit methods. We then use the result() method to acquire and display the results.

5.3 The Bell State

Renowned as the "Hello World" of Quantum Computation, the Bell State demonstrates two of the major concepts of Quantum Computing with the use of just two gates. By executing the Bell State circuit, we can observe both the Superposition and Entanglement of two qubits.

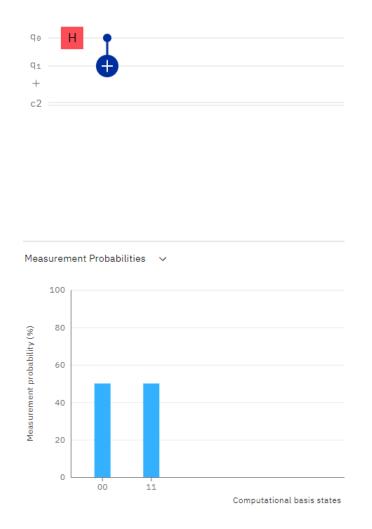


Figure 5-2 Bell State Experiment.

We can see from the measurement probabilities that we have two possible system states. Either both qubits will be zero or both qubits will be one. By applying a Hadamard gate on qubit 0 we transcend it to a state of equally zero and one and by applying a CX gate with the controller qubit being zero, whenever qubit zero is zero, qubit one will be zero and whenever qubit zero is one, qubit one will be one. Hence the two qubits are now entangled, which means that the state of qubit zero is now the major contributing factor of what the state of qubit one will be.

```
5.3.1 Implementation
```

5.5.1 Implementation
from qiskit import IBMQ
from qiskit import (
QuantumCircuit,
execute,
Aer)
class HelloWorld:
@classmethod
def run(cls):
<pre>with open('./credentials/token', 'r') as file:</pre>
<pre>token = file.read()</pre>
<pre>IBMQ.save_account(token, overwrite=True)</pre>
Use Aer's qasm_simulator
<pre>simulator = Aer.get_backend('qasm_simulator')</pre>
Create a Quantum Circuit acting on the q register
circuit = QuantumCircuit(2, 2)
Add a H gate on qubit 0
circuit.h(0)
Add a CX (CNOT) gate on control qubit 0 and target qubit 1
circuit.cx(0, 1)
Map the quantum measurement to the classical bits
circuit.measure([0, 1], [0, 1])
Execute the circuit on the qasm simulator
job = execute(circuit, simulator, shots=1000)
Grab results from the job
result = job.result()
Returns counts
<pre>counts = result.get_counts(circuit)</pre>
<pre>print("\nTotal count for 00 and 11 are:", counts)</pre>
Draw the circuit
print(circuit)

Figure 5-3 Code for Method run() of Experiment Bell State.

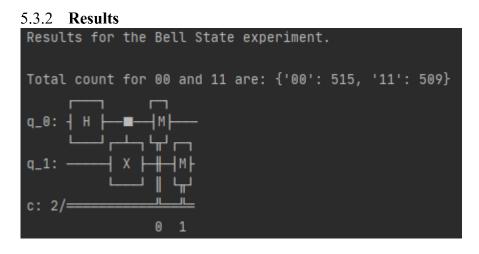
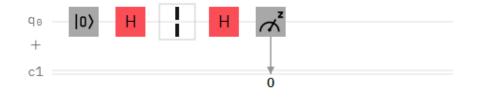


Figure 5-4 Results for Experiment Bell State.

We can indeed see that the only two possible outcomes are the states $|00\rangle$ and $|11\rangle$ with almost equal occurrences.

5.4 Phase Kickback

Phase Kickback is a phenomenon that occurs in Quantum Computation, and the reason why quantum probabilities are different from classical probabilities. If we were to imagine that a Hadamard gate's effect is similar to that of a coin flip, we would expect that two consecutive Hadamard gate operations on the same qubit would be analogous to a coin being flipped twice. Due to phase kickback, that is not the case.



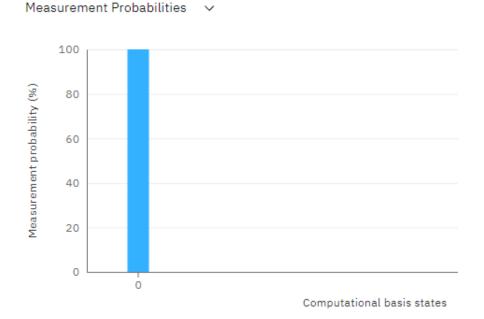


Figure 5-5 Phase Kickback effect.

Why does this happen? How can we have a definite outcome by applying randomness?

Assume a qubit that is instantiated in the $|0\rangle$ state.

Applying a Hadamard gate to this qubit will result in the following complex state:

$$H|0\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix} \cdot \begin{bmatrix} 1\\ 0 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1\\ 1 \end{bmatrix} = |+\rangle$$

The $|+\rangle$ state, when classically measured in the Z-basis, yields $|0\rangle$ or $|1\rangle$ with equal probability. But we can easily see that when we apply two Hadamard gates to the same qubit we have the following outcome:

$$H|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix} \cdot \frac{1}{\sqrt{2}} \begin{bmatrix} 1\\ 1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 1 \cdot 1 + 1 \cdot 1\\ 1 + (-1 \cdot 1) \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 2\\ 0 \end{bmatrix} = \frac{1}{2} \cdot 2 \begin{bmatrix} 1\\ 0 \end{bmatrix} = \begin{bmatrix} 1\\ 0 \end{bmatrix} = |0\rangle$$

We can see that after two consecutive applications of the Hadamard gate, the amplitudes of the mixed states cancel out and we are left with the definite state of $|0\rangle$.

We can see that if we choose certain linear combinations, some possible states may be destructed, or constructed.

Phase Kickback plays a vital role in Quantum Algorithm Development because it helps to make a transition from a seemingly random system to a useful result. We need to be careful in order to destruct the states that we are not interested in and construct the states that interest us.

5.4.1 Implementation

from qiskit import execute, Aer, QuantumCircuit
class Interference:
@classmethod
<pre>def run(cls):</pre>
Use Aer's qasm_simulator
<pre>simulator = Aer.get_backend('qasm_simulator')</pre>
Create a Quantum Circuit acting on the q register
circuit = QuantumCircuit(1, 1)
Add a H gate on qubit 0
circuit.h(0)
Add another H gate to qubit 0
circuit.h(0)
Map the quantum measurement to the classical bits
circuit.measure(0, 0)
Execute the circuit on the qasm simulator
<pre>job = execute(circuit, simulator, shots=1000)</pre>
Grab results from the job
result = job.result()
Returns counts
<pre>counts = result.get_counts(circuit)</pre>
<pre>print("\nTotal count for 0 and 1 are:", counts)</pre>
print(circuit)

Figure 5-6 Code for the run()Method of the Phase Kickback Experiment.

5.4.1.1 Results

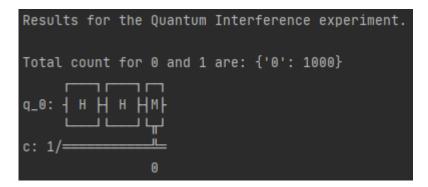


Figure 5-7 Results for Phase Kickback.

Indeed, we see that the only state measured is the $|0\rangle$ state.

This concludes the chapter. It is worth noting that for all these simple experiments we used the simulator in order to observe the behavior of an errorless system without further complications.

Chapter 6

QCLG Level 2 – The Exciting Game

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6.1 Concept

We have explored a plethora of different concepts in quantum computing and familiarized ourselves with some of the bizarre and unconventional attributes of quantum computation. As a comprehension exercise, we have implemented a quantum minigame, which showcases some of the basic quantum concepts we have learnt.

The game is played like a two-player card game.

In this case, the cards are quantum gates and our playing field is two qubits, on which our cards operate.

The goal of the game is to apply the quantum gates on your possession in the correct order to manage to transcend your qubit from the ground state of $|0\rangle$ to the excited state of $|1\rangle$ before a certain number of turns.

After that, a round is completed, and points are awarded to both players accordingly. Once the number of rounds finished, the player that was "excited" for the most rounds, wins.

Hence the name of the game.

6.2 Rules

The rules are very simple.

- 1. Finite limit of rounds and the player with the most winning rounds is the winner.
- 2. For each round, there is a countdown where the player is given the choice to draw the card from the deck and place it on their hand or they can choose to ignore the draw and stay with their current hand.
- 3. Once the countdown is finished, is player is required to place their cards in the order they want them to be applied on their qubit.
- 4. The playing field is evaluated and the player who possesses an excited qubit is awarded a point.
- 5. The hands and playing field are reset for each round.

6.3 Available Cards

In the current version of the Exciting Game, there are four different available cards.

- **THE HADAMARD GATE**. At this point, we are familiar with the Hadamard Gate. The player can use it to insert a randomness effect if they wish to test their luck or come up with tricks using quantum interference.
- THE X GATE. The X gate is a very powerful card in the Exciting Game because it can effectively win a round by flipping a qubit from |0⟩ to |1⟩.
- **THE CX GATE.** The CX gate provides a sabotage element to the game. It can flip the qubit of the enemy player.
- **THE RX GATE.** The RX gate has a similar effect to the Hadamard gate with a few key differences. It only acts on the X-axis and applying it twice in a row will result in a qubit state flip instead of negating the effect. For the purposes of the game, it can be thought of as a "half" X gate. It is up to the players to discover all of its usages.

6.4 Winning Hands

Below we provide some examples of winning hands in order to get a better grasp of the game.

Example 1

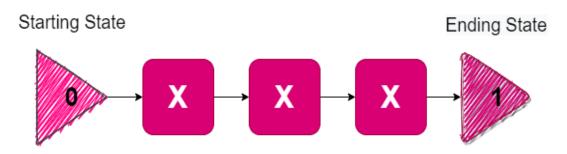
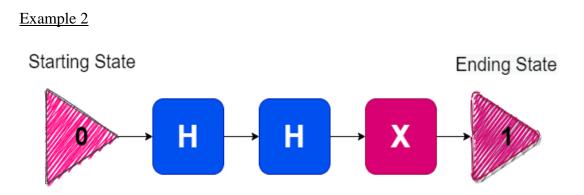


Figure 6-1 The Exciting Game Example 1.

In Figure 6-1 we can easily see how we get to the excited state. Applying the X gate for the first time flips the qubit from the $|0\rangle$ state to the $|1\rangle$ state. Applying the second X gate will lead to another state flip which will lead to the $|0\rangle$ once again. With the third application of the X gate, we end up with the $|1\rangle$ state which is a winning hand.



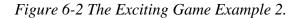


Figure 6-2 showcases a scenario where we can take advantage of interference in order to achieve a winning hand. By putting the two Hadamard gates one after the other we negate their effect and thus keeping intact the initial state of $|0\rangle$. With the application of a single X gate, we once again reach a winning position.

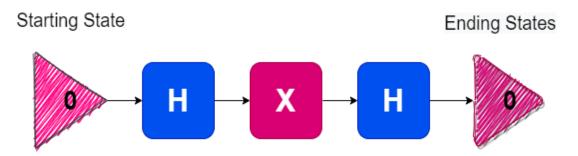


Figure 6-3 The Exciting Game Example 3.

In Figure 6-3, we have the same gates as in example 2. However, notice how the final state is $|0\rangle$ this time. By applying the X gate in the middle of the two Hadamard gates, due to the phase kickback effect, the X gate effect on the qubit is destroyed and we are left with the $|0\rangle$ state.

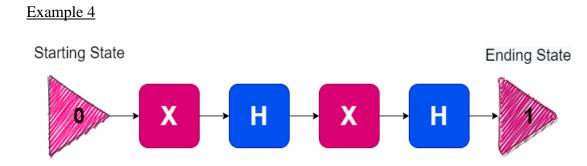
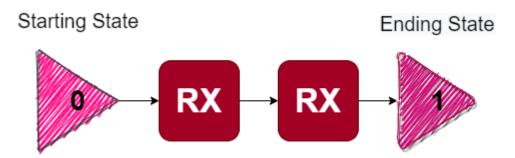
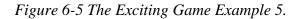


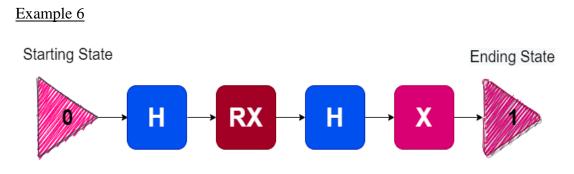
Figure 6-4 The Exciting Game Example 4.

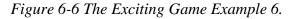
Building on the same logic from Figure 6-3, we can easily see why we reach the $|1\rangle$ state in this example. With the first application of the X gate, we reach the $|1\rangle$ state. Then, by surrounding the second X gate with two Hadamard gates we add nothing new to the $|1\rangle$ state, hence we measure the exciting state.





In Figure 6-5, we showcase the simplest winning position possible with two RX gates. With the first application of the RX gate, the qubit is rotated by $\pi/2$ in the X-axis and has an equal probability of measuring $|0\rangle$ or $|1\rangle$. However, with the second application, the qubit is rotated another $\pi/2$ in the X-axis hence formulating the absolute $|1\rangle$ state.





The RX gate is also subjectable to the phase kickback effect. In Figure 6-6, the first three gates effectively do nothing on the state of the qubits, and then by simply appending an X gate, we achieve the exciting state.



Figure 6-7 The Exciting Game Example 7.

Following the logic from Figure 6-4, we can see that with the first two applications of the RX Gate the qubit will be in the $|1\rangle$ state. With the application of the third and fourth RX Gate, the qubit will once again reach the $|0\rangle$ state. Hence, by finally appending the X gate we reach the desired state.

Example 8

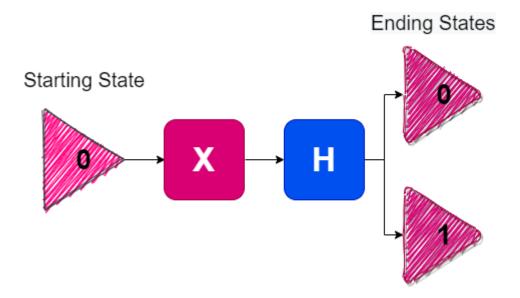


Figure 6-8 The Exciting Game Example 8.

By this point, it should be no surprise that in each round it is not always possible to achieve the excited state. This is an example of a hand where despite having an X gate, the Hadamard gate causes the qubit to behave with uncertainty. There is an element of luck in this case.

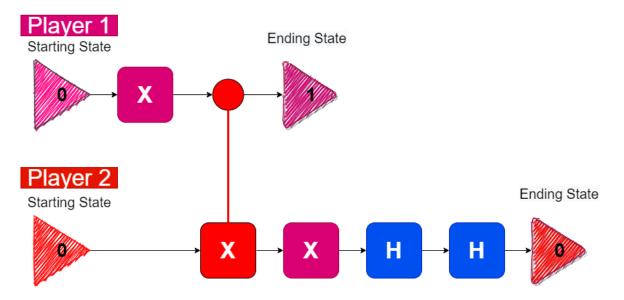


Figure 6-9 The Exciting Game Example 9.

The CX Gate is what makes the game interesting. We can see that in this scenario, player 1 applies an X gate hence putting their qubit to the $|1\rangle$ gate. As we know the CX gate is a Controlled Not gate so when player one applies a CX gate they will manage to flip the starting state of the player-two qubit. Now player 2's evaluation starts with the $|1\rangle$ state as the initial state. Hence, when player 2 applies their X gate they are now to the $|0\rangle$ state. Then with two Hadamard gates, nothing is achieved and player 2 ends up losing the round. We can also consider the scenario where player two had a CX gate in their possession. The best play for player 2 would be to insert their CX gate as their first card to flip the state of player 1 so that neither of them wins the round.

There are is a wide range of different combinations and outcomes that each round can follow with just four different gates. We have just shown a few. We have found a few good practices in terms of strategy.

6.5 Strategy

Through experimentation, or by simply playing the game, one can notice there are a few good practices when playing the Exciting Game.

Below we list a few strategies and a brief reasoning behind each strategy.

1. Drawing a card is not always a good idea.

In a surprising number of games, we have found that more often than not, you can construct a winning hand before all cards are dealt for the current round. You need to careful in order to realize that you possess a hand capable of winning even if you have two or three cards.

2. Phase kickback is your friend.

Even when a hand seems unpromising, there is almost always a way to destroy our less useful cards by using phase kickback. Consider a hand where the cards are Hadamard, Hadamard, X, X. At first glance, this is not a good hand since we have an even number of X gates. However, if we take advantage of phase kickback, we can put the one X gate in-between the two Hadamard gates and destroys its effect leaving us with effectively one X gate which is enough to make our qubit excited. You should always look for a possible phase kickback combo.

3. When all else fails, sabotage your opponent.

This might seem obvious but even if you can't reach the exciting state, you might still be able to sabotage your opponent by applying a CX gate at a point when your qubit is for sure excited.

6.6 Implementation

This is the driver program for the game. All is done in a while-loop until the number of rounds finishes.

Below we list some of the most useful methods implemented.

6.6.1 Method Definitions

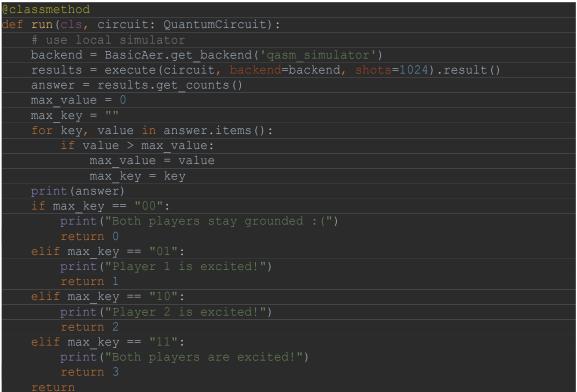


Figure 6-10 Code for the run() Method of The Exciting Game.

Method Definition – run(circuit):

The purpose of this method is to declare the winner after the end of each round.

Input: The circuit containing the two players' qubits along with their placed gates.

Output: A number between 0 and 3 indicating the four different outcomes.

Data Structures Used: In order to find the outcome, we need to traverse through the dictionary of measurements and find the state measurement with the most counts.



Figure 6-11 Code for place_gate() Method of The Exciting Game.

Method Definition – place_gate(player, field, qubit):

The purpose of this method is to place the player's gates onto the qubit in the correct order.

Input: The player's hand, the circuit containing the two players' qubits, the qubit of the current player.

Output: A modified field circuit.

Data Structures Used: In order to place the gates to the player's qubit we pop the items in the player list one by one and according to the card's name we append the corresponding quantum gate to the designated qubit of the field circuit, given by the qubit parameter. One thing to note is the way of assigning the CX gate. If are dealing with player one then that means that we want the CX gate to be placed with qubit 0 as the control qubit and qubit 1 as the target qubit and vice-versa.

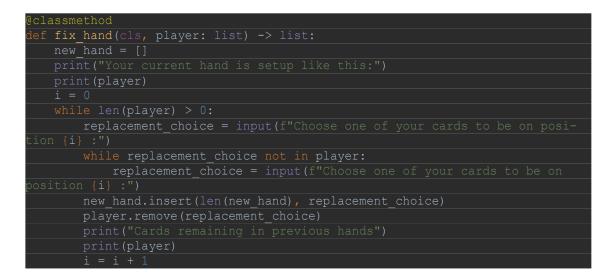


Figure 6-12 Code for the fix_hand() Method of The Exciting Game.

Method Definition –fix_hand(player):

The purpose of this method is to sort the player's gates according to their liking.

Input: The player's hand.

Output: The new, modified, hand of the player, as designated by the player's choices.

Data Structures Used: We display the current hand of the player. Then, in a while-loop we make the player choose the next card they want in their ordered final submission for the round. We do this by removing each of the valid choices of the player until their hand is exhausted and we have constructed a new list containing the previous cards of the player but in the order that they would like them to be executed on their qubit.

This concludes the implementation of the Exciting Game. The whole code for this project can be found in the appendix under the Exciting Game section.

Chapter 7

QCLG Level 3 – Implementation of Algorithms

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7.1 Deutsch - Jozsa

One of the first Quantum algorithms with a substantial speedup over its classical counterpart [5]. This is an algorithm that exploits the ability of a quantum system to support a huge state space in order to calculate the outcome faster than a conventional processor. For the problem this algorithm tackles, let f be a function that accepts as input either one or zero and outputs either one or zero. This is what we will call a Boolean function.

$$f: \{0,1\} \to \{0,1\}$$

Deutsch – Jozsa's task is to determine, given a function f, if it is given a Balanced Boolean function or a Constant Boolean function. Consider now the two Boolean types of functions: Balanced and Constant. A balanced function is a function where for all possible inputs, the outcome is zero for exactly half the possible inputs and one for the exact remaining half. A constant function is labelled as such when, regardless of input, its outputs are either only zero or only one.

> balanced: f(0) = 0, f(1) = 1balanced: f(0) = 1, f(1) = 0constant: f(0) = 1, f(1) = 1, constant: f(0) = 0, f(1) = 0,

Figure 7-1 Constant-Balanced Function Definition.

7.2 Classical Approach

Let's go through a scenario for a one-bit input for a classical environment.

We want to determine if a one-bit function f is balanced or constant given the fact that it can only be either balanced or constant and nothing else. We have two possible inputs. Zero and one. We will simply then query the function twice and after accumulating the results we will determine by aggregating the results if the function is balanced. In simpler words by receiving from the output a single zero and a single one or if the function is balanced by receiving two ones or two zeroes.

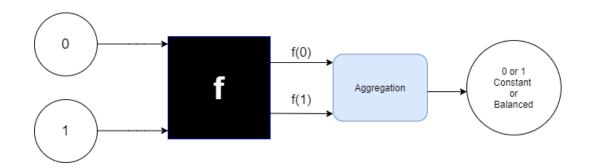


Figure 7-2 Deutsch-Jozsa Classical Approach Abstract Form.

Classical Deutsch-Jozsa Structure

Input: Sequences of bit-strings.

Processing: Based on the oracle function, map each input to an output bit

Output: 'Balanced' if outputted both zero and one **or** 'Constant' if outputted only zero or only one.

7.2.1 Classical Implementation

Below we provide the code for both the aggregation-oracle function and the driver function that calls the algorithm.

Class ClassicalXor:

This class contains three methods and imports DateTime for time measurement and the BitCombinations Class that contains helper methods for generating the worst-case scenario.

Method Definition - _super_secret_black_box_function_f(list_of_inputs):

This function is responsible for executing the full classical version of Deutsch-Jozsa. We provide a list of binary strings and then check all of them. Once we find both a one and a zero, we can say that it is a balanced function and if not, it is a constant. We assume that we are given either a balanced or constant function.

Input: An arbitrary sequence of bit strings.

Output: 'Balanced' or 'Boolean'.

Data Structures Used: The output_bit, which is used as a medium for executing multiple xor operations for each input. The output_zero and output_one variables are used as Boolean flags which are responsible for keeping track of what kind of values the output bit gets after each element. The counts variable keeps track of the amount of checks. The method is implemented by taking a list and then with a loop iterate over each element of the list. In an encapsulated loop, we take each element and execute a xor operation to the output bit with all bits of the element. We then check the output_bit value and change the value of the flags accordingly. We then reset it to zero and proceed with the conditional statements before continuing with the next iteration. The first conditional statement checks if both flags are set which means we have managed to find both possible values of the output bit hence we terminate and return a 'Balanced' answer. The second conditional statement checks if we have reached $2^{n-1} + 1$ inputs checked. If that is the case that means we still haven't managed to set both flags, hence our function is Constant.

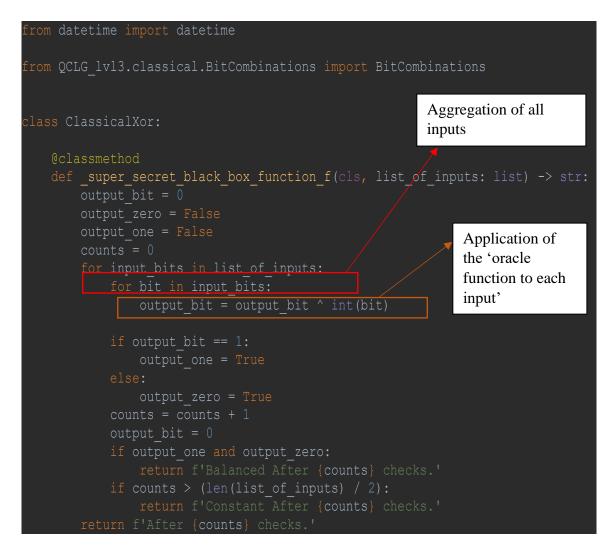


Figure 7-3 Code for Constructing the Classical Oracle for Deutsch-Jozsa.

Method Definition - execute_classical_xor(bits):

Input: Number of bits

Output: The nature of the function, 'Balanced' or 'Boolean', and the time it took for aggregating all inputs.

Data Structures Used: The bits parameter is an integer value that determines the length of the inputs that are going to be fed to the super_secret_black_box_function_f. We must first generate a list that contains all possible bit combinations. For this experiment, we have opted to always generate the worst-case scenario of inputs. We then declare two timestamps using the DateTime class for measuring the time it took for the input list to be generated and another two timestamps that measure the time it took our black box function evaluation to be finished. We then return in a list that contains two strings, one with the output of the black box function evaluator method we discussed above and one with the execution times.

```
@classmethod
def execute_classical_xor(cls, bits) -> list:
    start = datetime.now()
    inputs = BitCombinations.produce_worse_scenario(BitCombinations.combina-
tions(bits))
    end = datetime.now()
    elapsed = end - start
    time_to_generate_worst_input = elapsed.total_seconds()
    start = datetime.now()
    function_nature = cls._super_secret_black_box_function_f(inputs)
    end = datetime.now()
    elapsed = end - start
    time_str = elapsed.total_seconds()
    final = [time_to_generate_worst_input, time_str, bits, function_nature]
    return final
```

Figure 7-4 Code for Executing the Classical Solution of Deutsch-Jozsa.

By increasing the number of input bits to n we can see that the number of possible inputs increases exponentially and thus the number of queries to the function f. The reason we chose to pass as input the worst-case scenario every time is that several possible scenarios could occur and drastically change the performance of our classical solution. For example, if chose to randomly generate n-bit sequences, we could query the function f for the first time with our first input and get the answer one and then query it for the second time and get the answer zero. Given the fact we can only expect a balanced or constant function, we immediately know that by receiving both possible inputs, the function cannot be constant, hence it is balanced. These scenarios are mainly affected by the nature of the function -the way it reacts to each input- and by the way, the input is fed to the function. The latter can greatly increase or decrease the number of queries it takes to identify the function. Since the quantum systems that were available to us were not very powerful, we decided to always use the worst-case scenario in order to pursue a meaningful comparison. Let's say lay down the absolute worst scenario for verifying that a function is balanced. After an arbitrary way of feeding inputs to our function f we manage to force our function to output all its ones at once and all its zeroes, if there are any after the ones have finished.

7.2.2 Worst-case Scenario for a Classical Solution

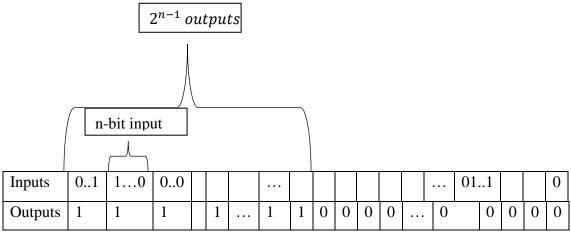


Figure 7-5 Worst-Case Scenario for Classical Deutsch-Jozsa.

It is easy to see that to verify the function is either constant or balanced we will have to query $2^{n-1} + 1$ times in the worst case in order to verify that the function is balanced or constant, in this case balanced.

Below we showcase three helper functions, the main function that generates all possible combinations and the function that groups the output of the main function in a way that simulates the worst-case scenario in the BitCombinations Class.

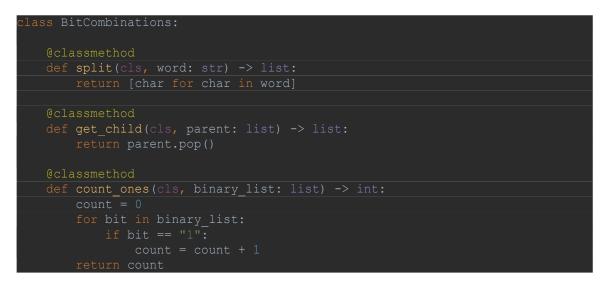


Figure 7-6 Helper Methods for Constructing the Worst-Case Scenario.

Method Definition - split(word): For translating a binary string into a list of '0' and '1'.

Input: String containing zeroes and ones.

Output: List containing zeroes and ones as characters.

Data Structures Used: The input parameter is a string word and the method returns a list representation of the string that contains all the characters of the string. We do that by iterating over each character of the list and by encapsulating that in square brackets we return a list of those characters.

Method Definition - get_child(parent):

Input: A list containing lists of binary strings

Output: The last element of the list.

Data Structures Used: This method takes a list by reference and pops and returns the last element of the list. The reason for this small method is to make the main method more readable and abstract.

Method Definition - count_ones(binary list): This function is used to count the number of ones contained in a current list. It is helpful when we want to check if we can add any more zeroes in the next iteration.

Input: A list containing ones and zeroes

Output: The number of ones.

Data Structures Used: The parameter is a list containing characters that are '0' or '1'. We initialize a count variable and we iterate through each character and increment whenever we find a '1' and return count.

Given a number of bits n, we want a method that will find all the possible combinations of n-bit numbers with k-bits set to one where $0 \le k \le n$. The solution should print all numbers with one set bit first, followed by numbers with two bits set, etc, up to the numbers whose all n-bits are set to one. The only cases where we manually add lists if for the all-zero and all one lists, which are added first and last. The number k symbolizes how many '1' we will have in each iteration.

Tree Diagram for 4-bit combinations.

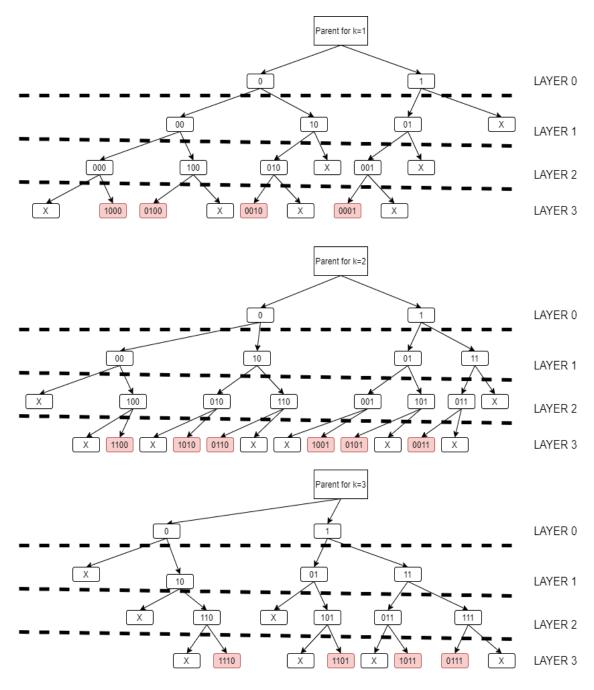


Figure 7-7 All different Combinations for 4 bits.

<pre>def combinations(bits: int) -> list:</pre>
final layer = []
for n bits in range(1, bits):
parent = []
for layer in range(bits - 1):
new layer = []
if layer == 0:
parent = []
b0 = ["0"]
b1 = ["1"]
parent.append(b0)
parent.append(b1)
<pre>while len(parent) > 0:</pre>
child = get_child(parent)
new_child1 = ["1"]
<pre>new_child0 = ["0"]</pre>
<pre>remaining_ones_for_child = n_bits - count_ones(child)</pre>
<pre>remaining_length = bits - len(child)</pre>
<pre>if remaining_ones_for_child == remaining_length:</pre>
<pre>new_child1.extend(child)</pre>
<pre>new_layer.append(new_child1)</pre>
<pre>elif 0 < remaining_ones_for_child < remaining_length:</pre>
new_child1.extend(child)
<pre>new_layer.append(new_child1)</pre>
<pre>new_child0.extend(child)</pre>
<pre>new_layer.append(new_child0)</pre>
<pre>elif remaining_ones_for_child == 0:</pre>
<pre>new_child0.extend(child)</pre>
<pre>new_layer.append(new_child0)</pre>
parent = new_layer
final_layer.append(parent)
result = []
<pre>all_zeros = [[split("0" * bits)]]</pre>
<pre>all_ones = [[split("1" * bits)]]</pre>
result.extend(all_zeros)
result.extend(final_layer)
result.extend(all_ones)
return result

Figure 7-8 Code for Generating n-bit Combinations.

Method Definition - Combinations(bits): This method is responsible to generate all bit combinations and return a list containing all of them.

Input: The number of bits n to determine the length of the n-bit combinations

Output: A list that contains all possible combinations with a number of '1' ranging from zero to n

Data Structures Used: We used an iterative approach to tackle this problem. We initialize an empty list, final_layer, which in the end will contain all 2^n combinations. We then start a loop that will iterate all K parent lists. In each of this loop's iteration, we want to finish producing all the combinations with K bits set to '1'. In the code, K is named n_bits. The loop is started with K equal to one and ends with K equal to n-1.

Now, to complete each parent, we initialize an empty list and start to fill each layer. The number of layers needed to complete a parent is always equal to n. For layer zero, which is the first layer, we initialize our 'tree' by appending the two possible elements, '0' and '1' to the parent list.

Now for each layer, using a while loop we take each child of the parent and count how many ones, we are allowed to still add to that child according to our current K. We do that by using the variable remaining ones for child, which calculates the remaining ones by subtracting the current amount of ones that the child list contains using the count ones() method we described previously and subtracting that from the n bits. We then calculate the remaining length of the child list by subtracting the current length of the child from the number of bits n. We are now left with two options. Either append '1' or '0' to the beginning of the list. The way we do that and prepare the child for the next layer is by creating two lists at the beginning of the while loop, one starting with '1' and one starting with '0', named new child1 and new child0. We then extend to those newly created lists with the child of the previous layer and append the new child to the new layer thus slowly creating the next layer. We have the option of appending the old child list to the new child1 or the new child0 or both, to determine to which new children the old child list will be appended and continue its evolution to the new layer, a child must fulfil some properties which we lay down in conditional statements. If the child's remaining ones that need to be added are equal to the number of bits that this child is allowed, then we can only add ones to this child and not zeroes. All these scenarios are shown in Figure 7-7. If the first conditional statement is not true, we check if the remaining ones for the child are more than zero and less than the remaining length that the current child can take. In this case, we can create two new children from the previous child, and we do this by extending both new child0 and new child1 with the previous child and appending them both to the next layer.

Lastly, for the conditions, if the remaining ones are neither equal to the number of bits the child is allowed nor more than zero, we are only allowed to extend the new_child0 list, which means appending a new child that contains an additional zero to the front.

Once we have done this for all children of the current parent list, we and emptied it, we are left with the new_layer list that contains all possible sequences with one more bit than the previous layer. We assign that new layer as our new parent list and continue until we have reached the final layer iteration and the last parent list pretty much contains all possible combinations and we assign it to the final layer list.

We then create a new list result and manually append all the combinations with zero ones in the sequence, which is only one combination, then all the other combinations we have generated, and then all the combinations with n ones in the sequence which is again only one combination. So, to summarize, the result list which will be returned is a list that contains lists. Each list element result contains, has all combinations of k-bits set to one.

This is depicted nicely in Figure 7-9 for four bits.

```
Enter Nomber of Dits. **
Combinations for K=0 bits set to '1'.
[['0', '0', '0', '0']]
Combinations for K=1 bits set to '1'.
[['0', '0', '0', '1'], ['0', '0', '1', '0'], ['1', '0', '0'], ['0', '1', '0', '0']]
Combinations for K=2 bits set to '1'.
[['0', '0', '1', '1'], ['1', '0', '0', '1'], ['0', '1', '0', '1'], ['1', '0', '1', '0'], ['0', '1', '1', '0'], ['1', '1', '0', '0']]
Combinations for K=3 bits set to '1'.
[['1', '0', '1', '1'], ['0', '1', '1'], ['1', '1', '0', '1'], ['1', '1', '0']]
Combinations for K=4 bits set to '1'.
[['1', '1', '1', '1']]
```

Figure 7-9 The n-bit Combinations Grouped by Number of Ones.

Now we need to re-sort the result list to have its first 2^{n-1} elements grouped by what they would make the black box function output. In the case of the xor operation, having an odd number of bits will output one, and having an even will output zero.

With that in mind, we simply must choose all k-one combinations that are either of an odd or even sum of ones, for our experiment to put odd sums first and group them together and only then append the even groups. Looking back to Figure 7-5 we can see that we have successfully simulated the worst possible scenario for the classical computer, which is feeding our algorithm all inputs that make the balanced function output one. Only then supply it with an input that will provide zero to the output, thus confirming a Boolean function We do this with balanced the help of the produce worse scenario method:

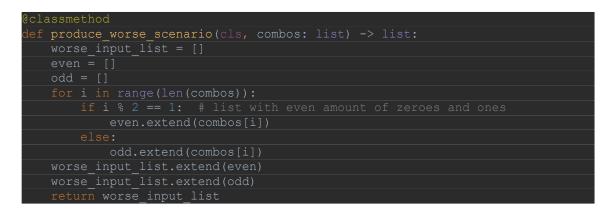


Figure 7-10 Code for Rearranging the bit Combinations to Construct Worst Scenario.

Method definition - produce_worse_scenario(combos): Take the list that contains all lists of lists of ones and zeroes and group them into odd and even groups. Then append all odd groups to one list and the even group afterwards to produces a list containing all the combinations of uneven amounts of ones and zeroes and combinations of ones and zeroes with the same number of ones and zeroes in a row following.

Input: A list containing lists, where each corresponding list contains all combinations for 1 to n-1 ones.

Output: A sorted list that groups even and uneven numbers of ones and zeroes.

Data Structures Used: The combos parameter is the result list from the BitCombinations method. We initialize an empty list worse_input_list and in a loop, we check to see if the current i modulo 2 is odd. We do that because we know how the structure of the combos list is. We know that for four bits for example when i is equal to zero, the list contains all the combinations with zero ones which when given to the black box function will be balanced. For i equal to one we have an odd number of ones, 1, for i equal to two we have two ones, etc. This is an easy way to group our lists.

Below we can see the worst possible input for four bits.

```
Worst input
['0', '0', '0', '1']
['0', '0', '1', '0']
['1', '0', '0', '0']
['0', '1', '0', '0']
['1', '0', '1', '1']
['0', '1', '1', '1']
['1', '1', '0', '1']
['1', '1', '1', '0']
['0', '0', '0', '0']
['0', '0', '1', '1']
['1', '0', '0', '1']
['0', '1', '0', '1']
['1', '0', '1', '0']
['0', '1', '1', '0']
     '1', '0', '0']
 '1', '1', '1', '1']
```

Figure 7-11 Output of Worst Scenario for 4 bits.

7.2.2.1 Results for Worst-case scenario

After applying the worst-case scenario, we always need to check for $2^{n-1} + 1$ times.

 $12 \ bits = 2^{12} = 4,096 \ and \ for \ 2^{n-1} + 1 \ checks \ we \ have \ 2^{11}$

+ 1 which is 2049 checks.

Figure 7-12 Execution Times for Classical Solution for 12 bits.

 $23 \ bits = 2^{23} = 8,388,608 \ and we have 2^{22} + 1 \ which is 4,194,305 \ checks.$

Enter:
1 for execution on the Local Device Simulator.
2 for execution on the Qasm simulator.
3 for execution on a real device.
4 for execution on the local device and real device!
Your input:1
Enter number of bits for a the classical solution:23
***************** FINAL * RESULTS ************************************
Balanced After 4194305 checks.
Determining if xor is balanced for 23 bits took 14.249531 seconds.
Time to generate worse input for 23 bits took 51.785056 seconds.

Figure 7-13 Execution times for Classical Solution for 23 bits.

24 *bits* = $2 * 2^{23} = 16,777,216$ *and* we have $2^{23} + 1$ which is 8,388,609 *checks*.

Enter:
1 for execution on the Local Device Simulator.
2 for execution on the Qasm simulator.
3 for execution on a real device.
4 for execution on the local device and real device!
Your input:1
Enter number of bits for a the classical solution:24
***************** FINAL * RESULTS ************************************
Balanced After 8388609 checks.
Determining if xor is balanced for 24 bits took 30.403439 seconds.
Time to generate worse input for 24 bits took 101.773189 seconds.

Figure 7-14 Execution Times for Classical Solution for 24 bits.

We can see that for each additional bit added the calculation time is exponentially increased. For just 32 bits it would take approximately 7 hours to calculate the worst-case scenario and approximately 2 hours to come to a solution.

7.3 Quantum Approach

We have set up a problem that is exponentially more complex. With the Deutsch – Jozsa algorithm we can devise a scheme for calculating all the possible outcomes in polynomial time instead of exponential using the quantum parallelism concept. Now we will work on an example of n = 2 qubits. Hadamard gates are used to prepare a superposition of all four possible states, 2^2 . We also need to implement the function f, or oracle function as it is called, in some way. When Deutsch – Jozsa is explained, the oracle function is often interpreted as a black box function that uses quantum parallelism to compute all values of f(x), x being one of the four possible configurations in our example.

$$|\alpha\rangle = \frac{1}{2}|00\rangle + \frac{1}{2}|01\rangle + \frac{1}{2}|10\rangle + \frac{1}{2}|11\rangle$$

Another concept that makes Deutsch -Jozsa work, is interference, which is achieved with the help of Hadamard gates once again. After applying them for a second time to the values that the oracle function outputted the measurement will be zero for all constant functions and one for balanced functions.

In our experiment, we made use of the oracle function implemented with CNOT gates, which are the quantum equivalent of the traditional XOR gate.

XOR Truth ta	ıble
00	0
01	1
10	1
11	0

CNOT Truth	table
00	0
01	1
10	1
11	0

Table 1 XOR and CNOT Truth Tables.

7.3.1 Algorithm's Structure

If we were to descript an abstract form of the algorithm it would comprise four major steps.

Step 1: Prepare the superposition of the qubits.

Step 2: The oracle function f.

Step 3: Interference.

Step 4: Measurements.

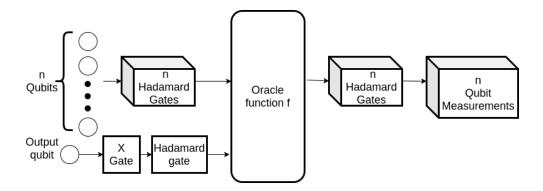


Figure 7-15 Deutsch Jozsa Structure.

This is the general form of the Deutsch-Jozsa for n qubits. For n qubits, we prepare a superpositioned state of 2^{n+1} different states because of the addition of the output qubit. This state is passed into the oracle function and all possible states are evaluated in polynomial time. We then proceed with the decoupling of this quantum state by taking advantage of Hadamard being its inverse and phase kickback.

7.3.2 **Proof of Concept with a worked example**

Using Dirac notation as a tool we can prove that in an errorless system we can expect to get a consistent output for both types of functions. For this worked example, we have two-qubit input and must show that for a balanced function the output bits measure one and for a constant, zero. It is worth noting that the output qubit is not measured in the final step since we can effectively and reliably determine the function while only measuring the input qubits.

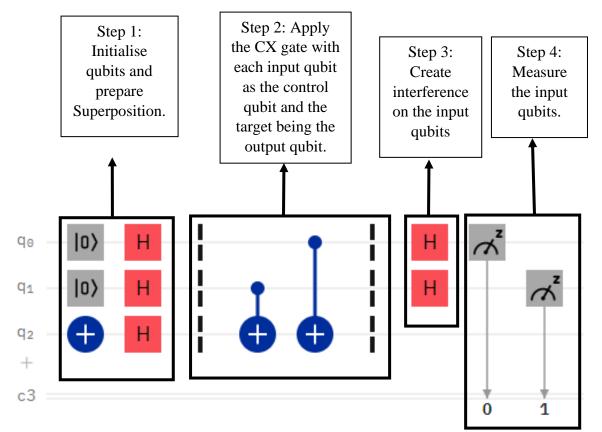


Figure 7-16 Worked Example.

7.3.2.1 Step 1 – Initialization

We initialize two-qubit registers. One contains 2 qubits, q0, and q1, and is initialized to the $|0\rangle$ state and the other one contains one qubit, q2, which is initialized to the $|1\rangle$ state.

$$|\psi_0\rangle = |0_0 0_1\rangle * |1_2\rangle$$

We then create a superposition of the three qubits in the following manner:

$$|\psi_1\rangle = \frac{1}{\sqrt{2}}(|0_0\rangle + |1_0\rangle) * \frac{1}{\sqrt{2}}(|0_1\rangle + |1_1\rangle) * \frac{1}{\sqrt{2}}(|0_2\rangle - |1_2\rangle)$$

Using the distributive property:

$$|\psi_1\rangle = \frac{1}{2}(|0_00_1\rangle + |0_01_1\rangle + |1_00_1\rangle + |1_01_1\rangle) * \frac{1}{\sqrt{2}}(|0_2\rangle - |1_2\rangle)$$

$$\begin{split} |\psi_1\rangle &= \frac{1}{2\sqrt{2}} (|0_0 0_1 0_2\rangle + |0_0 1_1 0_2\rangle + |1_0 0_1 0_2\rangle + |1_0 1_1 0_2\rangle - |0_0 0_1 1_2\rangle - |0_0 1_1 1_2\rangle - \\ |1_0 0_1 1_2\rangle - |1_0 1_1 1_2\rangle) \end{split}$$

7.3.2.2 Step 2 - Oracle

We have created a superposition of all possible states and we now pass this 'variable' into our function. Our function, which is essentially two C-NOT gates with the input qubits as the control qubits and the output qubit as the target qubit, maps the state $|x\rangle|y\rangle$ to the state $|x\rangle|y \oplus f(x)$). In order to understand this mapping let's consider the C-NOT gate's properties. It does not alter the control qubit's state. Hence $|x\rangle$ is outputted as itself. Control qubits q0 and q1 follow the $|x\rangle$ mapping. However, the $|y\rangle$ qubit of the function, in our case the output-target qubit q2 will have its value XORed with the outcome of the next CNOT operation on itself. So, for our case, qubits zero and one will remain the same, but qubit two will be transformed according to our oracle's implementation in a manner the following manner:

From $|\psi_1\rangle$ in Step 1, after passing through the oracle function, the system will be in this state:

$$\begin{split} |\psi_{2}\rangle &= \frac{1}{2\sqrt{2}} \Big(|0_{0}0_{1} \left((0_{2} \oplus f(0_{0})) \oplus f(0_{1}) \right)_{2} \rangle + |0_{0}1_{1} \left((0_{2} \oplus f(0_{0})) \oplus f(1_{1}) \right)_{2} \Big) + \\ |1_{0}0_{1} \left((0_{2} \oplus f(1_{0})) \oplus f(0_{1}) \right)_{2} \Big) + |1_{0}1_{1} \left((0_{2} \oplus f(1_{0})) \oplus f(1_{1}) \right)_{2} \Big) \\ f(1_{1}) \Big)_{2} \Big) - |0_{0}0_{1} \left((1_{2} \oplus f(0_{0})) \oplus f(0_{1}) \right)_{2} \Big) - |0_{0}1_{1} \left((1_{2} \oplus f(0_{0})) \oplus f(1_{1}) \right)_{2} \Big) - \\ |1_{0}0_{1} \left((1_{2} \oplus f(1_{0})) \oplus f(0_{1}) \right)_{2} \Big) - |1_{0}1_{1} \left((1_{2} \oplus f(1_{0})) \oplus f(1_{1}) \right)_{2} \Big) \Big) \end{split}$$

As we know, $x \oplus 0 = x$ and $x \oplus 1 = \overline{x}$ so $|\psi_2\rangle$ is simplified to:

$$\begin{split} |\psi_{2}\rangle &= \frac{1}{2\sqrt{2}} \Big(|0_{0}0_{1}(f(0_{0}) \oplus (0_{1}))_{2}\rangle + |0_{0}1_{1}(f(0_{0}) \oplus f(1_{1}))_{2}\rangle + |1_{0}0_{1}(f(1_{0}) \oplus f(1_{0}))_{2}\rangle + |1_{0}0_{1}(f(1_{0}) \oplus f(1_{1}))_{2}\rangle - |0_{0}0_{1}(\overline{f}(0_{0}) \oplus f(0_{1}))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0_{0}) \oplus f(1_{1}))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0_{0}) \oplus f(1_{1}))_{2}\rangle - |0_{0}1_{1}(\overline{f}(1_{0}) \oplus f(1_{1}))_{2}\rangle - |1_{0}0_{1}(\overline{f}(1_{0}) \oplus f(0_{1}))_{2}\rangle - |1_{0}1_{1}(\overline{f}(1_{0}) \oplus f(1_{1}))_{2}\rangle \Big) \end{split}$$

Before we proceed to Step 3, we need to consider the nature of our function. As discussed above, we can summarize by saying that a constant function is where f(0) = f(1) and a balanced function is where $f(0) = \overline{f}(1)$

If Constant:

$$\stackrel{|\psi_{2}\rangle}{\Longrightarrow} |\psi_{c}\rangle = \frac{1}{2\sqrt{2}} \Big(|0_{0}0_{1}(f(0) \oplus f(0))_{2}\rangle + |0_{0}1_{1}(f(0) \oplus f(0))_{2}\rangle + |1_{0}0_{1}(f(0) \oplus f(0))_{2}\rangle + |1_{0}0_{1}(f(0) \oplus f(0))_{2}\rangle + |1_{0}0_{1}(f(0) \oplus f(0))_{2}\rangle - |0_{0}0_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle - |1_{0}0_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle - |1_{0}0_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle \Big)$$

After factoring:

$$|\psi_{c}\rangle = \frac{1}{2\sqrt{2}} (|0_{0}0_{1}\rangle + |0_{0}1_{1}\rangle + |1_{0}0_{1}\rangle + |1_{0}1_{1}\rangle) * ((f(0) \oplus f(0))_{2} - (\overline{f}(0) \oplus f(0))_{2})$$

Simplifying the contents of qubit zero by replacing $f(x) \oplus f(x)$ with zero and $\overline{f}(x) \oplus f(x)$ with one.

$$|\psi_{c}\rangle = \frac{1}{\sqrt{2}}(|0_{0}\rangle + |1_{0}\rangle) * \frac{1}{\sqrt{2}}(|0_{1}\rangle + |1_{1}\rangle) * \frac{1}{\sqrt{2}}(|0_{2}\rangle - |1_{2}\rangle)$$

If Balanced:

$$\begin{split} \stackrel{|\psi_{2}\rangle}{\Longrightarrow} |\psi_{B}\rangle &= \frac{1}{2\sqrt{2}} \Big(|0_{0}0_{1}(f(0) \oplus f(0))_{2}\rangle + |0_{0}1_{1}(f(0) \oplus \overline{f}(0))_{2}\rangle + |1_{0}0_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle + |1_{0}0_{1}(\overline{f}(0) \oplus \overline{f}(0))_{2}\rangle + |1_{0}0_{1}(\overline{f}(0) \oplus \overline{f}(0))_{2}\rangle - |0_{0}0_{1}(\overline{f}(0) \oplus f(0))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0) \oplus \overline{f}(0))_{2}\rangle - |0_{0}1_{1}(\overline{f}(0) \oplus \overline{f}(0))_{2}\rangle - |1_{0}0_{1}(f(0) \oplus \overline{f}(0))_{2}\rangle - |1_{0}0_{1}(f(0) \oplus \overline{f}(0))_{2}\rangle \Big)$$

Simplifying the contents of qubit zero by replacing $f(x) \oplus f(x)$ with zero and $\overline{f}(x) \oplus f(x)$ with one.

$$\begin{split} |\psi_B\rangle &= \frac{1}{2\sqrt{2}} (|0_0 0_1 0)_2\rangle + |0_0 1_1 1_2\rangle + |1_0 0_1 1_2\rangle + |1_0 1_1 0_2\rangle - |0_0 0_1 1_2\rangle - |0_0 1_1 0_2\rangle - \\ |1_0 0_1 0_2\rangle - |1_0 1_1 1_2\rangle) \end{split}$$

After factoring:

$$|\psi_B\rangle = \frac{1}{2\sqrt{2}}(|0_00_1\rangle - |0_01_1\rangle - |1_00_1\rangle + |1_01_1\rangle) * (|0\rangle - |1\rangle_2)$$

$$|\psi_B\rangle = \frac{1}{\sqrt{2}}(|0_0\rangle - |1_0\rangle) * \frac{1}{\sqrt{2}}(|0_1\rangle - |1_1\rangle) * \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle_2)$$

7.3.2.3 Step 3 - Interference

Interference is the process of applying a Hadamard gate to each input qubit.

If Constant:

$$\stackrel{|\psi_{c}\rangle}{\Longrightarrow} |\psi_{cI}\rangle = \frac{1}{\sqrt{2}} (|0_0\rangle + |1_0\rangle)^H * \frac{1}{\sqrt{2}} (|0_0\rangle + |1_0\rangle)^H * \frac{1}{\sqrt{2}} (|0_2\rangle - |1_2\rangle)$$

$$|\psi_{cI}\rangle = |0_0\rangle * |0_1\rangle * \frac{1}{\sqrt{2}} (|0_2\rangle - |1_2\rangle)$$

If Balanced:

$$\stackrel{|\psi_B\rangle}{\Longrightarrow} |\psi_{BI}\rangle = \frac{1}{\sqrt{2}} (|0_0\rangle - |1_0\rangle)^H * \frac{1}{\sqrt{2}} (|0_0\rangle - |1_0\rangle)^H * \frac{1}{\sqrt{2}} (|0_2\rangle - |1_2\rangle)$$
$$|\psi_{BI}\rangle = |1_0\rangle * |1_1\rangle * \frac{1}{\sqrt{2}} (|0_2\rangle - |1_2\rangle)$$

We can see that after interference, the input qubits are in the ground state if the function is constant and in the excited state if the function is balanced. The driving force between all this trickery is the Hadamard gate, which as we have previously found out, is its inverse. Matrix multiplication clears the situation quite easily.

$$\begin{split} |\psi_{H}\rangle &= \left(\frac{1}{\sqrt{2}}(|0_{0}\rangle - |1_{0}\rangle)^{H} & \Leftrightarrow & |\psi_{H}\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \\ |\psi_{H}\rangle &= \frac{1}{2} \begin{bmatrix} 1*1+1*(-1) \\ 1*1+(-1)(-1) \end{bmatrix} & \Leftrightarrow & |\psi_{H}\rangle = \frac{1}{2} \begin{bmatrix} 0 \\ 2 \end{bmatrix} \\ |\psi_{H}\rangle &= \frac{1}{2}*2* \begin{bmatrix} 0 \\ 1 \end{bmatrix} & \Leftrightarrow & |\psi_{H}\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} & \Leftrightarrow \\ |\psi_{H}\rangle &= |1\rangle \end{split}$$

Likewise, if the input qubit is in the state,

$$\begin{split} |\psi_{H}\rangle &= (\frac{1}{\sqrt{2}}(|0_{0}\rangle + |1_{0}\rangle)^{H} \qquad \Leftrightarrow \qquad |\psi_{H}\rangle = \begin{bmatrix} 1\\ 0 \end{bmatrix} \qquad \Leftrightarrow \\ |\psi_{H}\rangle &= |0\rangle \end{split}$$

7.3.2.4 Step 4 - Measurements

We now apply two measurements to qubit zero and qubit one. If the function is constant, we expect two zeroes and if the function is balanced, we expect two ones. Since we know that f-CNOT is a balanced function we expect to see one from both qubits.



Figure 7-17 Measurement Probabilities for Balanced Function.

From Figure four we see that qubit zero's value is stored in bit zero in the classical bit register and qubit one's value is stored in bit one. We see that indeed one hundred percent of the time the result is that bot input qubits have the value one which indicates that the f-CNOT is a balanced function.

This concludes the logic behind the Deutsch-Jozsa algorithm and why it works. We now devise a plan to exploit this algorithm to determine a function's nature and compare it to a classical computer using the mainstream tools available to the modern software developer.

7.3.3 Quantum Implementation

The code is very straightforward. We have a class called DeutschJosza in which we will create our quantum circuit object.

Method definition - run_deutsch_josza(bit_string, eval_mode):

We first initialize a quantum circuit with n+1 qubits since we need to account for the output qubit as well and apply the necessary Hadamard gates to achieve a superpositioned state going into the oracle function appendment. We then apply interference and measure all input qubits.

Input: A bit sequence, a Boolean value dictating whether or not the algorithm will operate on eval mode, which essentially means not to print any display messages.

Output: The QuantumCircuit object.

Data Structures Used: In step one we create an empty QuantumCircuit object consisting of n+1 qubits and n bits. We then apply Hadamard gates to all input qubits using the h() method of the QuantumCircuit class and in a similar manner append an X gate and then a Hadamard gate to the output qubit.

In step two we create an oracle which is essentially another QuantumCircuit object. We do this with the CnotOracle Class. The create_cnot_oracle() method is responsible for returning an oracle of appropriate dimensions that works with CNOT gates.

Method Definition - create_cnot_oracle(input_string, input_length, eval_mode):

Input: The user-given input string, its length, and the operation mode

Output: A QuantumCircuit object

Data Structures Used: We create a QuantumCircuit object of input_length+1 qubits. The first and third for loops are there to initialize the qubits in a different state but in our experiments, we always start with the ground state of all qubits being zero. The second loop creates a CNOT operation of each input qubit with the output qubit. Before and after appending the CNOT operations we use two barrier objects. If the method is on eval_mode, then the printing messages will not be executed.

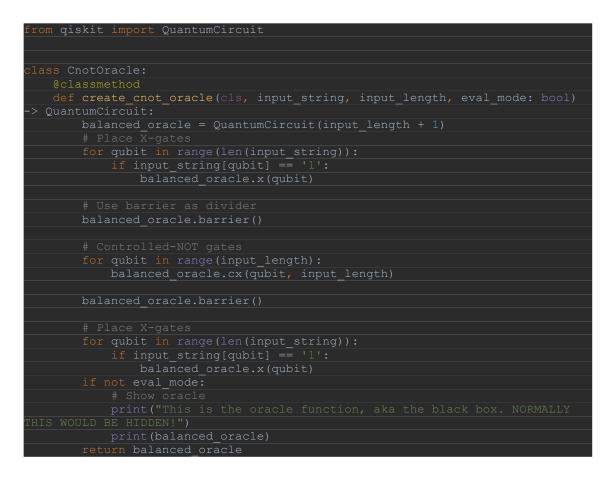


Figure 7-18 Code for Creating a Balanced Oracle.

Ins step 3 we apply Hadamard gates to all input qubits using a loop and then apply an additional barrier before applying measurements from all input qubits to the corresponding measurement bits.

The DeutschJosza class:

from qiskit import QuantumCircuit		
<pre>from oracles.cnot_oracle import CnotOrac class DeutschJosza:</pre>	ele 🔶	Step 1: Initialisation
<pre>@classmethod def deutsch_josza(cls, bit_string: s cuit:</pre>	tr, eval_mode: boc	ol) -> QuantumCir-
<pre>dj_circuit = QuantumCircuit(n + # Apply H-gates for qubit in range(n): dj_circuit.h(qubit) # Put output qubit in state -> dj_circuit.x(n) dj_circuit.h(n)</pre>	1, n)	Step 2: Oracle
<pre># Construct balanced oracle balanced_oracle = CnotOracle.cre eval_mode)</pre>	ate_cnot_oracle(bi	t_string, n,
<pre># Add oracle dj_circuit += balanced_oracle</pre>		
<pre># Repeat H-gates for qubit in range(n):</pre>	Ster Inte	o 3: rference
<pre># Measure for i in range(n): dj_circuit.measure(i, i) if not eval_mode: print(dj_circuit)</pre>	Step Meas	4: surements
<pre># return circuit return dj_circuit</pre>		

Figure 7-19 Code for Implementing Deutsch-Jozsa/

Indicative Results:

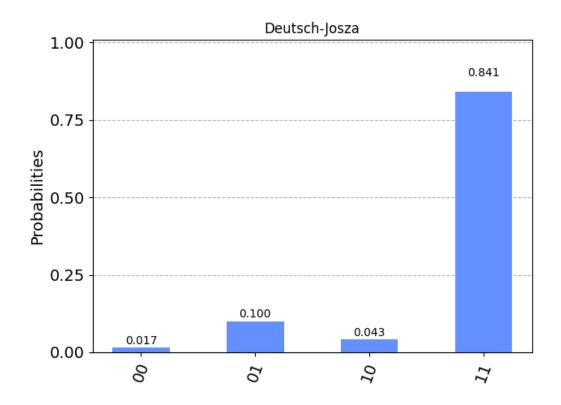


Figure 7-20 Different States Measured for Deutsch-Jozsa.

******************** FINAL * RESULTS ************************************
Total counts are: {'00': 17, '01': 102, '10': 44, '11': 861}
The time it took for the experiment to complete after validation was 5.900763034820557 seconds

Results of classical implementation for the Deutsch-Josza Algorithm:
Function is Balanced After 3 checks.
Time to generate worse input for 2 bits took 0.0 seconds.
Determining if xor is balanced for 2 bits took 0.0 seconds.

Figure 7-21 Results of a Classical and Quantum Execution for Deutsch-Jozsa.

We can see that indeed Deutsch-Jozsa outputs the "11" state which dictates that the C-NOT operation is balanced for many executions.

Also, we can see that our classical approach found the answer in 3 attempts, which for 2 bits goes according to the scheme we devised for the worst input which is $2^{n-1} + 1$.

7.4 Bernstein - Vazirani

Suppose a secret binary number X. We want to guess that number as quickly as possible, only knowing the length of the binary. We want to create a function that operates as few times as possible in order to find the secret number, [3].

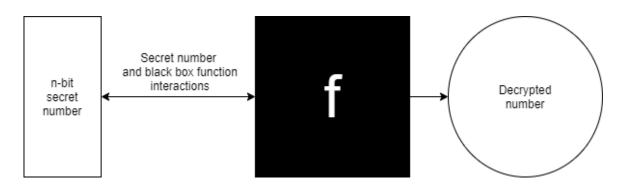


Figure 7-22 Bernstein-Vazirani Abstract Form.

How would a classical computer go to guess the number? One naïve approach would be to guess binary numbers of that length at random until we find it. It's easy to see that for a relatively long binary this would be rather hopeless since there are 2^n possible guesses. This is far from decreasing the number of possible iterations of our algorithm.

7.5 Classical Approach

There is of course a better way. We can take advantage of masking in order to guess an n-bit number in exactly n attempts each time. This is the exact algorithm below.

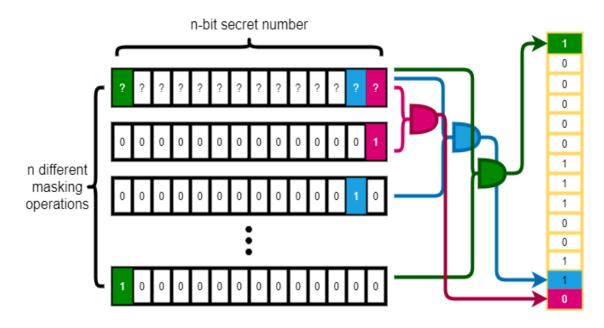


Figure 7-23 Bernstein-Vazirani Classical Approach.

We create n different n-length vectors containing only zero except in one place. We then proceed to take each vector and perform an AND operation with the secret number.

This way we decrypt a bit one by one. We start with the first vector making bit zero have the value one and all other n-1 bits have the value zero. If the zeroth bit is one then the AND operation will return one, which is the correct guess for the value of the secret bit, and if it returns zero then that is again a correct guess since zero AND one is equal to zero. All other bits of the vector will return zero since anything ANDed with zero returns zero. Hence, we can effectively check for each one of the secret bits and build the secret number in n guesses.

7.5.1 Classical Implementation Class BernsteinVaziraniClassical:

This class contains a single method that guesses the secret number. It may seem unnecessary to proceed with the masking operations since we already receive the secret number as a parameter. This is just for demonstration purposes. In a more realistic scenario, we would not have access to the actual secret number and the masking operation would be necessary since we would send our vector to be ANDed with the secret number, and the vendor which holds the secret number would return to us just the answer of the AND operation. Then this is the procedure we would follow. Since we will generate the secret numbers though, it rightly seems a bit counterintuitive.

Method Definition – guess_number(secret_binary):

Input: A string consisting of '1's and '0's.

Output: A string displaying the guess for the number and how many attempts it took.

Operate the AND operation on each individual bit of the secret number.

In python we can do this very easily without using a vector just by choosing which of the secret number to AND with 1. If the secret number was held by another person, we could do this with the same logic just by asking them to AND the ith bit of the number with 1 and tell us the result.

1

class BersteinVaziraniClassical:
Qclassmethod
<pre>def guess number(cls, secret binary):</pre>
mask = 1
guess = ""
attempts = 0
for bit in secret binary:
hit = int(bit) & mask
guess += str(hit)
attempts += 1
<pre>return f"My guess After {attempts} attempts is:\n{guess}."</pre>

Figure 7-24 Code for Guessing a Secret Number.

7.6 Quantum Approach

The Quantum approach is identical to that of Deutsch-Jozsa, we take advantage of interference while possessing the correct oracle function.

7.6.1 Algorithms' Structure

If we were to descript an abstract form of the algorithm it would again comprise of four major steps.

Step 1: Prepare the superposition of the qubits.

Step 2: The oracle function f.

Step 3: Interference.

Step 4: Measurements.

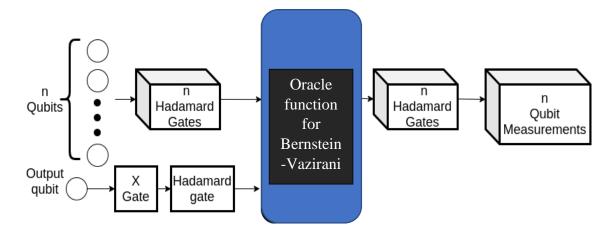


Figure 7-25 Bernstein-Vazirani Quantum Structure.

This is the general form of the Bernstein-Vazirani algorithm for n qubits. For n qubits, we prepare a superpositioned state of 2^{n+1} different states because of the addition of the output qubit. This state is passed into the oracle function and all possible states are evaluated with the oracle function. We then proceed with the decoupling of this quantum state by taking advantage of Hadamard being its inverse and phase kickback. As we can see, the structure of Bernstein-Vazirani is identical to that of Deutsch-Jozsa. The only real difference is the oracle function.

7.6.2 **Proof of Concept with a worked example**

Using Dirac notation as a tool we can prove that in an errorless system we can expect to get a consistent output for both types of functions. For this worked example, we have a two-bit number so we will need two input qubits and one output qubit and must show that we can successfully guess the number that is inside. It is worth noting that the output qubit is not measured in the final step since we can effectively and reliably determine the number while only measuring the input qubits. We will try and guess the secret number "01".

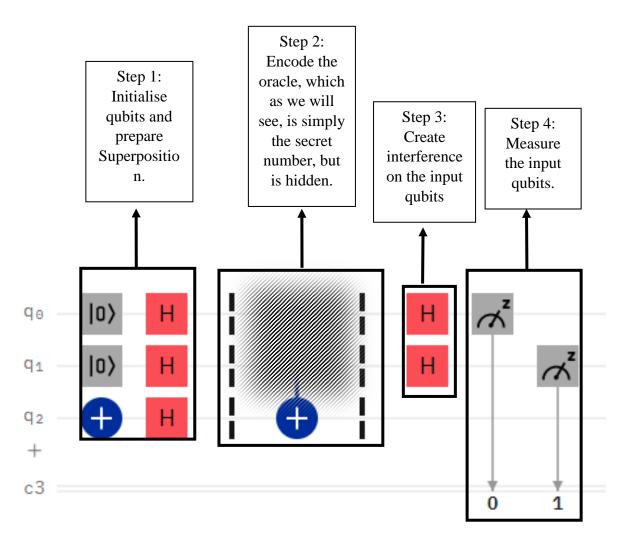


Figure 7-26 Bernstein-Vazirani Worked Example.

7.6.2.1 Step 1 – Initialization

We initialize our system the exact way we did in the Deutsch-Jozsa example.

We end up with the following system:

$$\begin{split} |\psi_1\rangle &= \frac{1}{2\sqrt{2}} (|0_0 0_1 0_2\rangle + |0_0 1_1 0_2\rangle + |1_0 0_1 0_2\rangle + |1_0 1_1 0_2\rangle - |0_0 0_1 1_2\rangle - |0_0 1_1 1_2\rangle - \\ |1_0 0_1 1_2\rangle - |1_0 1_1 1_2\rangle) \end{split}$$

7.6.2.2 Step 2 - Oracle

Our function, which is essentially a C-NOT gate wherever necessary to shape the secret number. The state $|x\rangle|y\rangle$ is once again mapped to the state $|x\rangle|y \oplus f(x)\rangle$. Like in Deutsch-Jozsa, qubits zero and one will remain the same, but qubit two will be transformed according to our oracle's implementation in a manner that we show below:

From $|\psi_1\rangle$ in Step 1, after passing through the oracle function, the system will be in this state:

$$\begin{split} |\psi_{2}\rangle &= \frac{1}{2\sqrt{2}} \Big(|0_{0}0_{1}(0_{2} \oplus f(0_{0}))_{2}\rangle + |0_{0}1_{1}(0_{2} \oplus f(0_{0}))_{2}\rangle + |1_{0}0_{1}(0_{2} \oplus f(1_{0}))_{2}\rangle + \\ |1_{0}1_{1}(0_{2} \oplus f(1_{0})))_{2}\rangle - |0_{0}0_{1}(1_{2} \oplus f(0_{0}))_{2}\rangle - |0_{0}1_{1}(1_{2} \oplus f(0_{0}))_{2}\rangle - |1_{0}0_{1}(1_{2} \oplus f(0_{0}))_{2}\rangle - |1_{0}0_{1}(1_{2} \oplus f(0_{0}))_{2}\rangle - |1_{0}0_{1}(1_{2} \oplus f(0_{0}))_{2}\rangle - |1_{0}0_{1}(0_{2} \oplus f(0_{0}))_{2}\rangle - |1_{0}0_{1}(0_{2} \oplus f(0_{0}))_{2}\rangle - |0_{0}0_{1}(0_{2} \oplus f(0_{0}))_{2}\rangle - |0_{0}0_{1}(0_{1} \oplus f(0_{0}))_{2}\rangle - |0_{0}0_{1}(0_{$$

As we know, $x \oplus 0 = x$ and $x \oplus 1 = \overline{x}$ so $|\psi_2\rangle$ is simplified to:

$$\begin{split} |\psi_{2}\rangle &= \frac{1}{2\sqrt{2}} \left(|0_{0}0_{1}f(0_{0})_{2}\rangle + |0_{0}1_{1}f(0_{0})_{2}\rangle + |1_{0}0_{1}f(1_{0})_{2}\rangle + \\ |1_{0}1_{1}f(1_{0})_{2}\rangle - |0_{0}0_{1}\overline{f}(0_{0})_{2}\rangle - |0_{0}1_{1}\overline{f}(0_{0})_{2}\rangle - |1_{0}0_{1}\overline{f}(1_{0})_{2}\rangle - |1_{0}1_{1}\overline{f}(1_{0})_{2}\rangle \right) \end{split}$$

Before we proceed to Step 3, we need to consider the nature of our function. The function f that we use is the CNOT operation which as we now know is balanced.

Hence $f(0) = \overline{f}(1)$ and $\overline{f}(0) = f(1)$.

The equation is transformed as follows:

$$\begin{split} |\psi_{2}\rangle &= \frac{1}{2\sqrt{2}} \left(|0_{0}0_{1}\overline{f}(1)_{2}\rangle + |0_{0}1_{1}\overline{f}(1)_{2}\rangle + |1_{0}0_{1}f(1)_{2}\rangle + \\ |1_{0}1_{1}f(1)_{2}\rangle - |0_{0}0_{1}f(1)_{2}\rangle - |0_{0}1_{1}f(1)_{2}\rangle - |1_{0}0_{1}\overline{f}(1)_{2}\rangle - |1_{0}1_{1}\overline{f}(1)_{2}\rangle \right) \end{split}$$

After factoring:

$$|\psi_{2}\rangle = \frac{1}{2\sqrt{2}}(|0_{0}0_{1}\rangle + |0_{0}1_{1}\rangle - |1_{0}0_{1}\rangle - |1_{0}1_{1}\rangle) * (\overline{f}(1) - f(1))_{2}$$

and once more,

$$|\psi_{2}\rangle = \frac{1}{2\sqrt{2}}(|0_{0}\rangle - |1_{0}\rangle) * (|0_{1}\rangle + |1_{1}\rangle) * (\overline{f}(1) - f(1))_{2}$$

7.6.2.3 Step 3 - Interference

Interference is the process of applying a Hadamard gate to each input qubit, and we can also ignore the third qubit now.

$$|\psi_3\rangle = \frac{1}{\sqrt{2}}(|0_0\rangle - |1_0\rangle)^H * \frac{1}{\sqrt{2}}(|0_1\rangle + |1_1\rangle)^H$$

After interference, the input qubits are excited if there was a CX operation performed on the qubit inside the oracle function, or in the ground state if there was not a CX operation in the corresponding qubit inside the oracle.

Hence our resulting state is as follows:

 $|\psi_3\rangle = |0\rangle * |1\rangle * |state of qubit 2\rangle$

7.6.2.4 Step 4 - Measurements

We then simply measure the first two input qubits and can expect with certainty to measure the secret number which was indeed 01.

7.6.3 **Quantum Implementation**

The code is very straightforward. We have a class called BernsteinVazirani in which we will create our quantum circuit object.

Method definition – bernstein_vazirani (random_binary, eval_mode):

Input: A given random binary number and a Boolean value determining if the method will run on eval_mode, that is not print any display messages.

Output: A QuantumCircuit object

Data Structures Used: This method, like in Deutsch-Jozsa is broken down into four abstract points.

In step one we create a QuantumCircuit object of input_length+1 qubits. Using a for loop, we append Hadamard gates to the input qubits. We then append an X gate and a Hadamard gate immediately after to the output qubit, thus completing the initialization step.

In step 2, we first construct the secret number oracle by calling the create_secret_number_oracle() method from the SecretNumberOracle class and appending it to the dj_circuit which is the primary circuit object.

In step 3, we need to implement interference. We do that by instantiating a new circuit object and with a for loop append Hadamard gate to the first n qubits. We then also append a barrier to emphasize the distinct steps of the algorithm, but it is not necessary.

In step 4, using a for loop we assign a measurement to be performed on each of the input qubits and the value of that measurement to correspond to the correct bit in the bits register. The BersteinVazirani class:

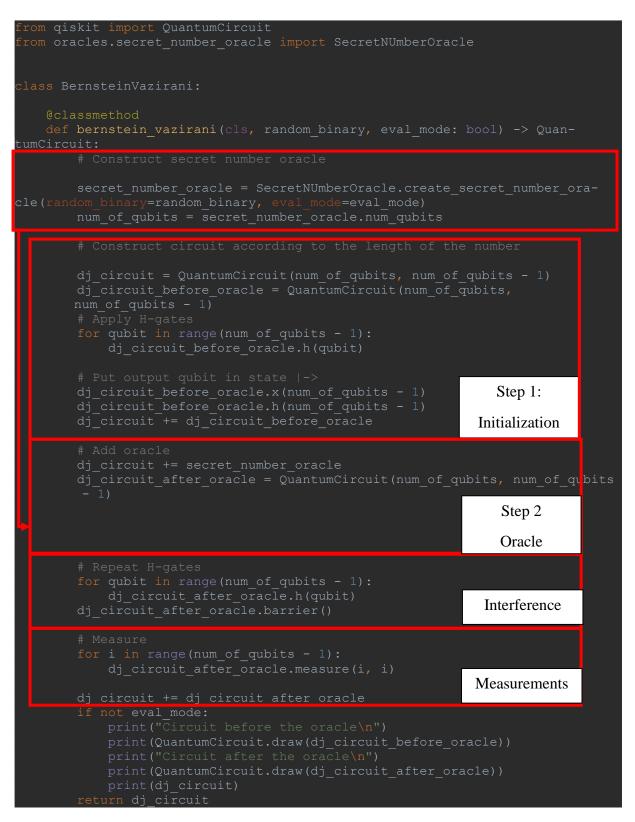


Figure 7-27 Code for Bernstein-Vazirani Implementation.

The SecretNumberOracle class:

This class is responsible for preparing the oracle function for the algorithm.

from qiskit import QuantumCircuit
class SecretNUmberOracle:
Qclassmethod
<pre>def create_secret_number_oracle(cls, random_binary, eval_mode: bool) -></pre>
QuantumCircuit:
n = len(random_binary)
<pre>secret_number_oracle = QuantumCircuit(len(random_binary) + 1,</pre>
len(random_binary))
Use barrier as divider
<pre>secret_number_oracle.barrier()</pre>
Controlled-NOT gates
<pre>for qubit in range(len(random_binary)):</pre>
<pre>if random_binary[qubit] == '1':</pre>
<pre>secret_number_oracle.cx(qubit, n)</pre>
<pre>secret_number_oracle.barrier()</pre>
if not eval_mode:
Show oracle
print("This is the oracle function, aka the black box. NORMALLY
THIS WOULD BE HIDDEN!")
<pre>print(secret_number_oracle)</pre>
return secret_number_oracle

Figure 7-28 Code for Creating the Oracle for the Bernstein-Vazirani Algorithm.

Method definition - create_secret_number_oracle(random_binary, eval_mode):

Input: A random binary and a Boolean value determining the method operation mode.

Output: A QuantumCircuit object

Data Structures Used: We create a QuantumCircuit object that has one more qubit than the length of the secret random binary. In the for-loop, we append a CNOT operation with the control qubit being one of the qubits that are in a position corresponding to that where one exists in the secret random binary. And the target is the output qubit. We then append once again a barrier to create a conceptual box around the oracle and return the circuit in order to complete the algorithm.

Indicative results:

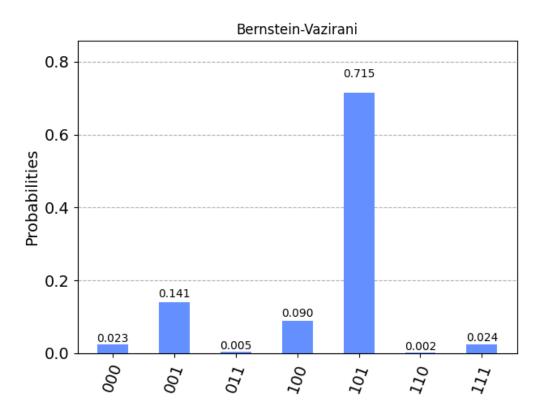


Figure 7-29 State Measurements for Bernstein-Vazirani.



Figure 7-30 Execution Results for Bernstein-Vazirani.

The secret number generated is 101. We can see that for the quantum implementation the correct state was measured for many executions.

The classical solution as we can see takes three guesses to find 101, which is a 3-bit number.

7.7 Summary

It is important that we stress the fact that both Deutsch-Jozsa and Bernstein-Vazirani are algorithms that solve black-box problems. Shor's Algorithms is another, more famous algorithm, that uses this black-box logic in order to factor integers in polynomial time [12].

Another major Algorithm, not covered in this thesis, is the Grover Search Algorithm, which is used to solve unstructured search problems using a phenomenon called amplitude amplification [8], which is not the exact same as Phase-Kickback. We quickly realize that Quantum Algorithm Development is highly dependent on Quantum Systems behavior and not algorithmic thought alone.

However, we need to realize, that, like Classical Computation has its own families of Algorithms, so does Quantum. There is a certain approach to tackling each different problem, and a good quantum intuition, is when we are able to recognise the nature of a problem in order to use the correct algorithm, much like conventional computing.

This concludes QCLG Level 3 / Implementation chapter.

Chapter 8

QCLG Level 4 – Evaluation of Algorithms

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8.1 Evaluation Process

Since quantum computing is not available as a local device in our home computer yet, we need to resort to cloud services. This means that we need to establish a connection with a real Quantum Computer and send our code to be executed in a real Quantum Machine and fetch back the results for evaluation. We also execute the classical solutions of the problems locally to be able to compare the execution times and discuss our findings. For this, we have implemented the Evaluation class which is responsible for executing both approaches multiple times in an automated manner and then displaying the results in a meaningful way. Some keynotes:

- We compare execution times for the classical execution for 1 bit and the quantum execution for 1 qubit, then for 2 bits and 2 qubits, etc.
- Each classical and quantum experiment is executed 1024 times for each different input size. This is done for two reasons. 1024 times is the default number of executions for a quantum experiment and we can observe more useful results since the execution times for executing each experiment once are very small it would be very difficult to extract useful information.
- For Deutsch-Jozsa specifically:
 - \circ We use a balanced oracle which we explained in section 7.3.2.2.
 - For the classical implementation, we again use the worst-case scenario as explained in section 7.2.2.
- For Bernstein-Vazirani specifically:
 - For the classical implementation, we used the sound classical method, which was explained here in section

After choosing the algorithm for evaluating the process is split as follows:

- 1. Classical Execution
- 2. Quantum Execution
- 3. Comparison

8.2 Implementation

We explain some of the important methods in the evaluation process.

<pre>from QCLG_lvl3.classical.random_binary import RandomBinary from tools import Tools</pre>
import constants
from qiskit.providers import JobStatus
class Evaluation:
@classmethod
<pre>def evaluate(cls, algorithm):</pre>
if algorithm == "0":
cls.evaluate deutsch josza()
elif algorithm == "1":
cls.evaluate_bernstein_vazirani()

Figure 8-1 Code for Redirecting to Algorithm Evaluation.

Method definition – evaluate(algorithm):

Input: Algorithm key

Output: Call to the corresponding evaluating method.

Data Structures Used: We check with an if, elif statement whether we are going to evaluate Deutsch-Jozsa or Bernstein Vazirani.

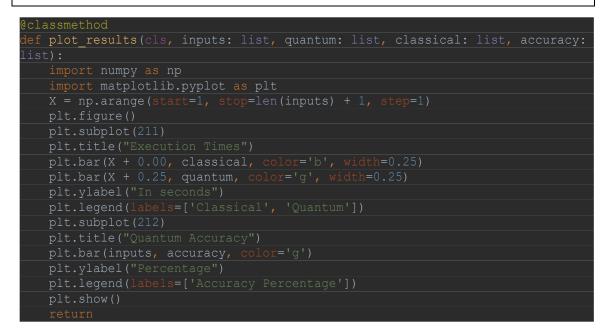


Figure 8-2 Code for Plotting Evaluation Results.

Method definition - plot_results(inputs, quantum, classical, accuracy):

A simple method for constructing two separate plots.

Input: The list containing the range of different inputs, the quantum time results, the classical time results, and the accuracy of the quantum results.

Output: A figure containing two different plots. The first plot is for displaying the execution times for both the classical and quantum executions for all different input sizes and the second plot is for displaying the percentage of correct answers the quantum system achieves for different input sizes.

Data Structures Used: For this method, we need to take advantage of the matplotlib capabilities. We instantiate a figure object and then we will create two subplot objects.

With the np.arange method we create a NumPy array that we will use for our X-axis numbers on both subplots. We will then define two new subplots using plt.subplot(211) and plt.subplot(212). The first two numbers in the parentheses mean that we are going to have two rows and one column of plots and then 1 for the first subplot and 2 for the second.

Deutsch-Josza Evaluation method

Method Definition – evaluate_deutsch_josza():

This method encapsulates the abstract process we are going to follow for evaluation. We will first generate all executions for the classical solutions of the algorithm, then proceed with the quantum solutions and finally plot the results to compare the two approaches effectively. This method is also greatly assisted by methods implemented in the Tools class which can be found at the end of this chapter.

Input: No input is required

Output: The figure containing the comparisons.

Data Structures Used: Using a while-loop we ensure that the input is a number between 1 and 14 since we are limited by the hardware available by IBM Q. We then define five lists. The circuits list will store the different quantum circuits that will be constructed for each different number of qubits. The n_bits list stores the number of qubits for each iteration. The quantum execution times list will store all the execution times for different each input size for the quantum circuits. The classical execution times list will store all the execution times for each different input size for the classical solutions. We then use a for-loop that will execute the classical solution for each size of the input range and append the execution time on the classical execution times list. We then display a completion percentage which is calculated easily by dividing the current iteration number, by the total number of iterations and multiplying by 100.

After we conclude with the classical executions, we need to execute the quantum solutions. This is done by finding the least busy backend capable of supporting the number of qubits required for the largest input we are going to evaluate and then submitting a batch job using the run_batch_job() from the Tools class and then using a while-loop we wait for all the circuits to finish executing in order to continue with the final assembly.

For each finished quantum experiment, we append the time to the quantum execution times list. We also retrieve the dictionary containing all the different measurements acquired from running each experiment 1024 times. We then search for the number of times the correct answer was measured. In the case of Deutsch-Jozsa and because we are running it with a balanced oracle, the correct answer is all input qubits measuring to 1. Hence, we are looking for the number of times the ('1' * amount of input qubits) is measured. If it was not measured, we assign the value 0. We then find the accuracy for the given input size by dividing by 1024, which is the time the experiment was executed for that input size and multiplying by 100. We append that percentage to the success rates list and continue with the rest of the input sizes. With the plotting conclusion of this while-loop, we call our method and the evaluate deutsch josza() method is completed.

The method definition for Bernstein-Vazirani is almost identical with the only difference being how to find how many times the correct answer was measured by the quantum system. To do this, we need to search in the dictionary of measurements the key which is equal to the secret binary number and find how many times that secret number was measured. We do this by storing the different random binaries in a separate list called random_binaries and then checking to see the count of the experiment for 3 qubits where the measurement is equal to the 3-bit secret number that was generated beforehand. The code for the Bernstein-Vazirani method can be found on Appendix D.

8.3 Deutsch-Jozsa Measurements

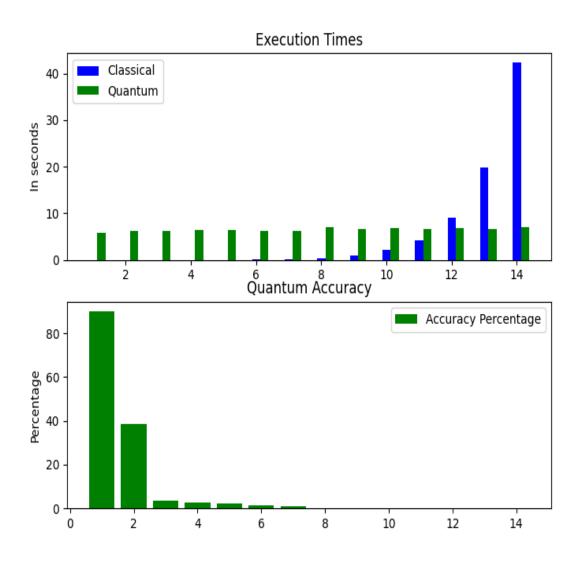


Figure 8-4 Deutsch-Jozsa Evaluation Measurements.

We can extract some very important information from the above figure. In the first subplot, we observe the execution times for input sizes of 1 to 14. We can see that for an input size of up to 11 bits/qubits, the classical computer is much faster. However, as we discussed in 7.1, Deutsch-Jozsa is an exponential problem. The exponential increase in time taken is obvious for input sizes 12, 13, and 14 with execution times increasing from 12 seconds to 24 seconds to 44 seconds respectively. At the same time, the execution time for the quantum system remains almost constant, a phenomenon that agrees with the fact that the quantum solution finds the answer with "one guess".

However, we cannot yet celebrate. While the quantum computer shows a lot of promise, timewise, we can see that we run into a lot of problems as the input size increases, accuracy wise.

Due to physical errors in quantum systems, the amount of times the quantum system measures the correct state is greatly decreased as we advance into greater input sizes. We can maybe empirically recognize the correct answer for up to 6 qubits, but from that point on, the correct measurement was sometimes not even measured once.

Error correction is another major section of quantum computing as discussed in detail here [10] and we can see from our findings, why it is so important.

8.4 Bernstein-Vazirani Measurements

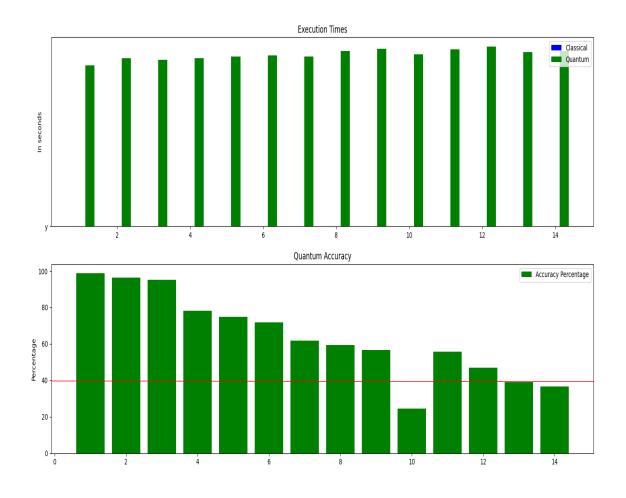


Figure 8-5 Bernstein-Vazirani Evaluation Measurements.

Continuing with the Bernstein-Vazirani algorithm, in the first subplot, we can see that the classical execution times are negligible even for plotting. That is because the Bernstein-Vazirani is solved in linear time in our classical solution and not in exponential time as Deutsch-Jozsa. The accuracy of Bernstein-Vazirani however, is much better. We can see that even for a 14-qubit secret number, the correct answer was measured almost 40% of the time which is substantially more accurate. We think that this is because of how the oracle function is implemented. Since Bernstein-Vazirani requires less C-NOT gates than Deutsch-Jozsa, there should be fewer physical errors, due to the smaller amount of gates.

8.5 Additional Assisting Methods

0.5	Additional Assisting methods
	# Evaluation methods
·	assmethod
def	<pre>prepare_dj(cls, bits: int):</pre>
	<pre>bit_sequence = "0" * bits</pre>
	<pre>dj_circuit = DeutschJosza.deutsch_josza(bit_sequence, eval_mode=True)</pre>
	return dj_circuit
@cla	assmethod
def	<pre>prepare bv(cls, random binary: str):</pre>
	bj circuit = BernsteinVazirani.bernstein vazirani(random binary,
evai	l mode=True)
	return bj circuit
0cla	assmethod
def	deutsch josza classical(cls, bits: int):
	return ClassicalXor.execute classical xor(bits=bits)
0cla	assmethod
def	bernstein vazirani classical(cls, bits: int):
	random binary = RandomBinary.generate random binary v2(bits)
	return BersteinVaziraniClassical.guess number(random binary)
0cla	assmethod
def	<pre>run batch job(cls, circuits: list, least busy backend) -> ManagedJobSet:</pre>
	transpiled circuits = transpile(circuits, backend=least busy backend)
	# Use Job Manager to break the circuits into multiple jobs.
	job manager = IBMQJobManager()
	job set eval = job manager.run(transpiled circuits,
bacl	kend=least busy backend, name='eval',
	max experiments per job=1) #
max	experiments per job =1 very important to get
	# individual execution times
	return job set eval

Figure 8-6 Code for Supplementary Methods of Evaluation.

The purpose of these methods is to use pre-existing methods from the Tools class but call them differently in order to automate some inputs and speed up the evaluation process. The methods prepare_dj(), prepare_bv(), deutsch_josza_classical(), are almost exact replicas with the slight difference that they run on eval mode, which just means that the printing messages are dismissed.

The bernstein_vazirani_classical() method differs only in the way of generating the random binary, which is by randomly generating a random binary number of length 'bits' directly and not by asking the user. The implementation for generate_random_binary_v2(bits) can be found in the appendix, along with the rest of the code.

Method definition - run_batch_job(circuits, least_busy_backend):

Input: A list of quantum circuits and the list busy backend capable of executing the largest circuit in the given list.

Output: A ManagedJobSet object that holds all the managed jobs/experiments along with all their information. Further information about the ManagedJobSet object can be found here [32].

Data Structures Used: We first need to use the transpile method which will transpile the experiments into a manageable batch job for the backend to execute in parallel. We then run this batch job which returns each finished experiment asynchronously. This means that it does not wait for all experiments to finish to return the results. This is why we created a while-loop which waited for all experiments to finish in the evaluate_deutsch_josza() and evaluate_bernstein_vazirani() methods previously. The max_experiments_per_job was a very important parameter for successfully evaluating the algorithms. By limiting the number of experiments per job to 1, we can have discrete execution times for each different input size. When we tried to run the batch job with the default settings, this was not possible, and we could not extract discrete execution times.

This concludes the QCLG Level 4 / Evaluation chapter.

Chapter 9

Conclusions

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9.1 **Conclusions**

Concluding this thesis, we first need to say that quantum computing is a marvelous attempt of humankind to push the limits of knowledge to unprecedented new grounds and we are thrilled that we could be a part of the Quantum revolution that is afoot. Most importantly, we completed our goals and can answer the four research questions that we beset at the beginning of this thesis in section 1.2.

Firstly, we completed our goals with the development of the "Quantum Computing Learning Gate". Through this platform, we managed to demonstrate how a software developer can take advantage of the abundant resources available and initiate their own quantum journey. We have demonstrated the most useful examples that helped us proceed with more advanced concepts on the matter. We have developed a quantum game that capitalizes on the capabilities of a quantum system in order to provide a joyful and interactive experience whilst keeping a consistent learning environment. We have completed a case study of algorithms on quantum and classical systems to showcase an additional point of view of understanding in the emerging scene of quantum research. We have implemented a practical workflow for evaluating algorithms using IBM Q's systems. Finally, we have gathered useful material that can be a jumping-off point for any aspiring quantum developers to initiate their journey and contribute in their own way.

Secondly, we managed to find answers to our three questions.

The first question was about who should be interested in quantum computing. Should it be reserved for researchers? Is it a subject that should be made compulsory in all universities across the world? The answer to both of these questions is no. Quantum computing is moving forward through a joint effort of a passionate community that is its borders are way beyond the scientific community. We are in the age of quantum computing and significant contributions have been made through the power of open-source projects. However, it is our personal belief that it should not be a compulsory course in computer science degrees. It differs in many ways from traditional computation, and its impact, although significant is not yet at the stage of refinement that is required to be taught to undergraduate students. We would advocate the notion of it being available as an elective course to more undergraduate or post-graduate curriculums since in 10 years it could change the ways we approach machine learning, optimization problems, and security.

The second question was is regarding the choice between the available resources online for initiating quantum endeavors. The answer to this is not absolute. It largely depends on the goals of the individual or organization. We, as students have chosen Qiskit since it is a framework that puts education high in its list of priorities by offering a variety of tutorials through the Qiskit textbook, online videos, and even has a community on Slack for any Qiskit related issue. For professionals or organizations that are interested in quantum solutions, the decision is a lot tougher. Since quantum computing is offered through cloud services there are a lot of factors to consider. If they have worked with Amazon Web Services for conventional solutions previously, they might find it easier to make a transition to Amazon Braket. If, however, we are talking about a startup that wants to try out the new technology it might be best to explore other options that may offer cheaper rates.

The third and final question that we can answer is whether pursuing a career specialized in quantum computing is a viable career path. This question was the main inducement for completing this thesis. Quantum computing is a fascinating subject but when it comes up in conversations, even in professional software environments, especially in Cyprus, it is thought of as a strictly theoretical section of Computer Science and deemed as a purely academic career path. In our opinion, this is not the case. Although in its early industrial stages, quantum computing could offer solutions to companies, and why not, Cypriot companies. Maybe not in the next 3 to 5 years but definitely into the foreseeable future, a quantum computing specialization could be as valuable as a software engineering or cybersecurity specialization. The cloud is much closer than we realize and with it the quantum or hybrid solutions available. We are optimistic and for sure keeping the opportunity of working as a quantum developer high in the list of career paths.

9.2 Limitations

During the completion of this thesis, we encountered many obstacles. Our first obstacle was the absence of previous friction with quantum computing. Quantum computing is not a course that is offered in the standard curriculum or any of the elective courses at the University of Cyprus. Therefore, we had to start with very small steps which were time-consuming.

Another limitation was the quantum physics aspect of quantum computing. Although there have been many leaps with quantum frameworks, there is still a large benefit in having previous experience with complicated concepts such as the superposition, entanglement, amplitudes, and all other phenomena found in the world of quantum mechanics.

The third limitation had to do with the resources that were available to us and were twofold. IBM Q generously offers quantum systems for free but for this thesis, it was difficult to extract the data that we hoped for. Quantum systems are still error-prone, and while we were able to plan legitimate workflows for implementing and evaluate algorithms, we still could not achieve an advantage with the quantum systems because of the very low accuracy. The other issue regarding resources was the queue time required for executing experiments. We started working on this thesis in early January of 2020 and the traffic on the IBM Q systems was much lower than what it is today. At times, we needed to wait for hours to successfully execute jobs on the desired backends. This is of course a good thing as well because it means that more and more people are interested in quantum research, but it certainly slowed down the progress of experimentation.

An additional limitation was getting familiar with the Qiskit framework in a short amount of time. Qiskit is a very rich framework that is open-source and is three years into development. Its public GitHub repository has over 100 contributors [33]. A framework of the magnitude demands thorough research through documentation study and multiple trial and error experimentations. We continuously discovered more and more features available that were vital to the development of a more complete version of QCLG. We were only able to experience a glimpse of what Qiskit can offer through its different elements and hundreds of Python classes. This brings us neatly into our next section which is our aspirations for QCLG.

9.3 Future Work

Quantum computing, due to its difficulty, and studying demands, is challenging. This also leaves a lot of space for improvement in our current work. Concerning the QCLG there is a lot of ground to be covered. The four levels could be enhanced by additional experiments, additional algorithms for analysis, and more evaluation options. Also, there can be a fifth level, which will contain real-world applications of algorithms. Furthermore, "The Exciting Game" is a very promising project that could evolve into a standalone application, much like Quantum Tetris [35]. Our idea is that the current code can be modified to work as a backend-side application, and we can deploy a front-end graphical interface that takes the game to the next level in terms of interaction and playability. Another idea is the creation of a separate domain from GitHub that presents key points of QCLG and references material relative to quantum computing in order to host a portal capable of generating traffic. Finally, we believe that it would be a good idea to explore other quantum resources besides IBM and Qiskit to expand our horizons.

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Appendix A

QCLG Level 1 Code

import constants
from QCLG lvll.hello quantum world import HelloWorld
from QCLG lvl1.interference import Interference
from QCLG lvll.single qubit superposition import SingleQubitSuperposition
from QCLG lvl1.three qubits superposition import ThreeQubitSuperposition
class SimpleExperimentsManager:
@classmethod
def showcase(cls):
print("Hi, these are the available experiments")
<pre>for i in range(len(constants.experiments)):</pre>
<pre>print(f"{i}. {constants.experiments[i]}")</pre>
<pre>choice = input(f"Which experiment would you like to try from 0 to</pre>
{len(constants.experiments)-1}? ")
while choice not in constants.acceptable experiment inputs:
choice = input(f"Which experiment would you like to try from 0 to
(len(constants.experiments) - 1)? ")
if choice == "0":
HelloWorld.run()
elif choice == "1":
SingleQubitSuperposition.run()
elif choice == "2":
ThreeQubitSuperposition.run()
elif choice == "3":
Interference.run()

```
from qiskit import execute, Aer, QuantumCircuit

class Interference:
    @classmethod
    def run(cls):
        # Use Aer's qasm_simulator
        simulator = Aer.get_backend('qasm_simulator')
        # Create a Quantum Circuit acting on the q register
        circuit = QuantumCircuit(1, 1)
        # Add a H gate on qubit 0
        circuit.h(0)
        # Add another H gate to qubit 0
        circuit.h(0)
        # Map the quantum measurement to the classical bits
        circuit.measure(0, 0)
        # Execute the circuit on the qasm simulator
        job = execute(circuit, simulator, shots=1000)
        # Grab results from the job
        result = job.result()
        # Returns counts
        counts = result.get_counts(circuit)
        print("NnTotal count for 0 and 1 are:", counts)
        print(circuit)
```

```
itom qiskit import (
    QuantumCircuit,
    execute,
    Aer)

class HelloWorld:
    @classmethod
    def run(cls):
        with open('./credentials/token', 'r') as file:
            token = file.read()
    IBMQ.save_account(token, overwrite=True)
    # Use Aer's gasm_simulator
    simulator = Aer.get_backend('qasm_simulator')
    # Create a QuantumCircuit(2, 2)
    # Add a H gate on qubit 0
    circuit.h(0)
    # Add a CX (CNOT) gate on control qubit 0 and target qubit 1
    circuit.cx(0, 1)
    # Map the quantum measurement to the classical bits
    circuit.measure([0, 1], [0, 1])
    # Execute the circuit on the gasm simulator
    job = execute(circuit, simulator, shots=1024)
    # Grab results from the job
    result = job.result()
    # Returns counts
    counts = result.get_counts(circuit)
    print("NnTotal count for 00 and 11 are:", counts)
    # Draw the circuit
    #
```

```
from qiskit import execute, Aer, QuantumCircuit, IBMQ
class SingleQubitSuperposition:
    @classmethod
    def run(cls):
        # Use Aer's qasm_simulator
        simulator = Aer.get_backend('qasm_simulator')
        # Create a Quantum Circuit acting on the q register
        circuit = QuantumCircuit(1, 1)
        # Add a H gate on qubit 0
        circuit.h(0)
        # Map the quantum measurement to the classical bits
        circuit.measure(0, 0)
        # Execute the circuit on the qasm simulator
        job = execute(circuit, simulator, shots=1000)
        # Grab results from the job
        result = job.result()
        # Returns counts
        counts = result.get_counts(circuit)
        print("\nTotal count for 0 and 1 are:", counts)
        provider = IBMQ.load_account()
        backend = provider.backend.shots=1000)
        # Grab results from the job
        result = job.result()
        # Returns counts
        counts = result.get_counts(circuit)
        print("\nTotal count for 0 and 1 are:", counts)
        provider = IBMQ.load_account()
        backend = provider.backend.shots=1000)
        # Grab results from the job
        result = job.result()
        # Returns counts
        counts = result.get_counts(circuit)
        print("\nTotal count for 0 and 1 are:", counts)
        pront("\nTotal count for 0 and 1 are:", counts)
        # Durw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\nTotal count for 0 and 1 are:", counts)
        print("\nTotal count for 0 and 1 are:", counts)
        # Draw the circuit
        print("\
```

from qiskit import execute, Aer, QuantumCircuit, IBMQ
class ThreeQubitSuperposition:
@classmethod
<pre>def run(cls):</pre>
Use Aer's qasm_simulator
<pre>simulator = Aer.get_backend('qasm_simulator')</pre>
Create a Quantum Circuit acting on the q register
circuit = QuantumCircuit(3, 3)
Add a H gate on qubit 0
circuit.h(0)
circuit.h(1)
circuit.h(2)
Map the quantum measurement to the classical bits
<pre>for i in range(3):</pre>
circuit.measure(i, i)
Execute the circuit on the qasm simulator
job = execute(circuit, simulator, shots=1024)
Grab results from the job
result = job.result()
Returns counts
counts = result.get counts(circuit)
<pre>print("\nTotal count all possible states are:", counts)</pre>
provider = IBMQ.load account()
backend = provider.backends.ibmq valencia
Execute the circuit on a real device
<pre>job = execute(circuit, backend=backend, shots=1024)</pre>
Grab results from the job
result = job.result()
Returns counts
counts = result.get counts(circuit)
print("\nTotal count for all possible states are:", counts)
Draw the circuit
print(circuit)

All the above are the entirety of the code for QCLG Level 1. Here there are implementations for experimenting with superposition, interference and entanglement. All these experiments are managed by a standalone QCLG Level 1 manager class.

Appendix B

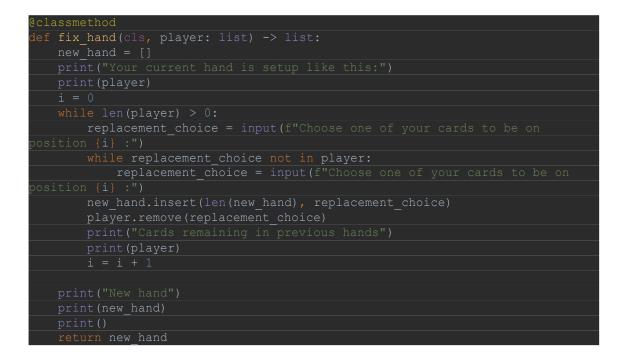
QCLG Level 2 Code

```
from random import randint
```

```
mport numpy as np
'rom qiskit import execute, BasicAer
'rom qiskit.circuit.quantumcircuit import QuantumCircuit
```

```
class Game
```

```
QuantumCircuit:
       cls.place gate(player2, field, 1)
   cards = ["H", "H", "X", "X", "CX", "RX", "RX"]
@classmethod
@classmethod
@classmethod
Qclassmethod
```



```
cls.shuffle deck(deck)
```

This code is for QCLG Level 2 and is the entire implementation of "The Exciting Game". There is a plethora of helper functions and the driver method play_the_exciting_game() which runs the game.

Appendix C

QCLG Level 3 Code

Classical Directory

class BersteinVaziraniClassical:
Qclassmethod
<pre>def guess_number(cls, secret_binary):</pre>
mask = 1
guess = ""
attempts = 0
<pre>for bit in secret_binary:</pre>
hit = int(bit) & mask
guess += str(hit)
attempts += 1
<pre>return f"My guess After {attempts} attempts is:\n{guess}."</pre>

```
mport random
class RandomBinary:
  @classmethod
  def generate_random_binary(cls, limit: int) -> str:
    rand = int(random.uniform(0, limit))
    print(f"random_binary in decimal:{rand}")
    random_bin = bin(rand)[2:]
    print(random_bin)
    return random_bin
  @classmethod
  def generate_random_binary_v2(cls, binary_length: int) -> str:
    random_bin = ""
    for i in range(binary_length):
        rand = int(random.uniform(0, 1))
        if rand >= 0.5:
            random_bin = random_bin + "1"
        else:
            random_bin = random_bin + "0"
    return random_bin
```

```
def _super_secret_black_box_function_f(cls, list_of_inputs: list) -> str:
BitCombinations.produce_worse_scenario(BitCombinations.combinations(bits))
function_nature]
```

```
@classmethod
def split(cls, word: str) -> list:
    return [char for char in word]
@classmethod
def get_child(cls, parent: list) -> list:
    return parent.pop()
@classmethod
def count_ones(cls, binary_list: list) -> int:
    count = 0
    for bit in binary_list:
        if bit == "1":
            count = count + 1
    return count
@classmethod
def produce_worse_scenario(cls, combos: list) -> list:
    worse_input_list = []
    even = []
    odd = []
    for i in range(len(combos)):
        if i % 2 == 1: # list with even amount of zeroes and ones
            even.extend(combos[i])
        else:
            odd.extend(combos[i])
    worse_input_list.extend(even)
    worse_input_list.extend(odd)
    return worse_input_list
```

```
from qiskit import QuantumCircuit

class CnotOracle:
    @classmethod
    def create_cnot_oracle(cls, input_string, input_length, eval_mode: bool)
    -> QuantumCircuit:
        balanced_oracle = QuantumCircuit(input_length + 1)
        # Flace X-gates
        for qubit in range(len(input_string)):
            if input_string[qubit] == '1':
                balanced_oracle.barrier()

        # Use barrier as divider
        balanced_oracle.cx(qubit, input_length)
        # Controlled-NOT gates
        for qubit in range(input_length):
            balanced_oracle.cx(qubit, input_length)
        balanced_oracle.barrier()

        # Place X-gates
        for qubit in range(len(input_string)):
            if input_string[qubit] == '1':
            balanced_oracle.cx(qubit, input_length)
        balanced_oracle.barrier()

        # Place X-gates
        for qubit in range(len(input_string)):
            if input_string[qubit] == '1':
                balanced_oracle.x(qubit)

        if not eval_mode:
            # Show oracle
            print("This is the oracle function, aka the black box. NORMALLY
THIS WOULD BE HIDDEN!")
        return balanced_oracle
        return balanced_oracle
```

```
from qiskit import QuantumCircuit

class SecretNUmberOracle:
    @classmethod
    def create_secret_number_oracle(cls, random_binary, eval_mode: bool) ->
QuantumCircuit:
        n = len(random_binary)
        secret_number_oracle = QuantumCircuit(len(random_binary) + 1,
len(random_binary))
        # Use barrier as divider
        secret_number_oracle.barrier()
        # Controlled-NOT gates
        for qubit in range(len(random_binary)):
            if random_binary[qubit] == '1':
                secret_number_oracle.barrier()
        if controlled-NOT gates
        for qubit in range(len(random_binary)):
            if random_binary[qubit] == '1':
                secret_number_oracle.barrier()
        if not eval_mode:
            # Show oracle
            print("This is the oracle function, aka the black box. NORMALLY
THIS WOULD BE HIDDEN!")
        print(secret_number_oracle)
        return secret_number_oracle
```

```
classical answer, algorithm):
   def showcase(cls):
```

The manager for handling algorithm execution choices. The different choices are to run the algorithm classically, on a real quantum device, on the IBM simulator or both on a quantum real device and classically on the local device.

```
QuantumCircuit:
eval mode)
```

```
QuantumCircuit:
SecretNUmberOracle.create secret number oracle(random binary=random binary,
num of qubits - 1)
```

Appendix D

QCLG Level 4 Code

<pre>from QCLG_lvl3.classical.random_binary import RandomBinary</pre>
from tools import Tools
import constants
from qiskit.providers import JobStatus
class Evaluation:
@classmethod
<pre>def evaluate(cls, algorithm):</pre>
<pre>if algorithm == "0":</pre>
cls.evaluate deutsch josza()
elif algorithm == "1":
cls.evaluate_bernstein_vazirani()
@classmethod
<pre>def plot results(cls, inputs: list, quantum: list, classical: list,</pre>
accuracy: list):
import numpy as np
import matplotlib.pyplot as plt
X = np.arange(start=1, stop=len(inputs) + 1, step=1)
plt.figure()
plt.subplot(211)
plt.title("Execution Times")
plt.bar(X + 0.00, classical, color='b', width=0.25)
plt.bar($X + 0.25$, quantum, color='q', width=0.25)
plt.ylabel("In seconds")
plt.legend(labels=['Classical', 'Quantum'])
plt.subplot(212)
plt.title("Quantum Accuracy")
plt.bar(inputs, accuracy, color='g')
plt.ylabel("Percentage")
<pre>plt.legend(labels=['Accuracy Percentage'])</pre>
plt.show()
return

Appendix D contains all code necessary for evaluating algorithms in level 4 of the QCLG.

retur

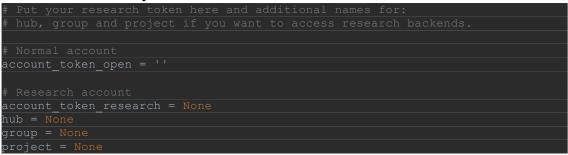
```
RandomBinary.generate random binary_v2(number_of_bits)
        completion_percentage = int((number_of_bits / test_range) * 100)
```

```
returr
```

Appendix E

Supplementary Code

Credentials Directory



Account_details.py. This file needs to be filled with the appropriate user information in order for the QCLG to function properly and be able to make calls to the remote IBM systems.

For this to happen the user needs to create an IBM Q account and generate their token on the starting page of login.

```
algorithms = ["Deutsch-Josza", "Bernstein-Vazirani"]
experiments = ["Bell State - Hello World", "Superposition with one Qubit",
acceptable_execution_inputs = ['0', '1', '2', '3', '4']
acceptable algorithm inputs = ['0', '1']
acceptable experiment inputs = ['0', '1', '2', '3']
cards = ["H", "H", "X", "X", "CX", "SX", "SXDG"]
acceptable choice inputs = ['1', '2', '3', '4']
experimentation level = "Enter:" \
```

Useful display messages.

Code for the manager class responsible for branching to the four levels of QCLG.

```
datetime):
        return difference.total seconds()
```

```
lef find least busy backend from open(cls, n):
tion().simulator and x.status().operational == True))
tion().simulator and x.status().operational == True))
@classmethod
rithm)]} on {least busy backend}")
```

```
answer = cls.execute bv in real device (random binary)
```

```
search(size)
rithms[int(algorithm)]} Algorithm:")
ate worst input} seconds.")
tion time} seconds.")
@classmethod
```

```
Evaluation methods
```

The supplementary code is required to support all actions involved in the QCLG levels. There is a plethora of backend calls to IBM resources, and calls to local implementations of classical solutions.

Appendix F

QCLG Setup

We demonstrate the procedure of setting up QCLG in a new Windows 10 machine. The instructions are very similar for Ubuntu with slight modifications required for paths. Clone the repository locally by doing the following:

- 1. Launch PyCharm.
- 2. From the "VCS" tab, choose the option "Get from Version Control..."
- 3. From the pop-up window paste the https URL of the dmsl/quantum repository: https://github.com/dmsl/quantum.git
- 4. Press the Clone button and choose to open on a new window.

The project will not yet work. We need to create and setup the conda environment and setup a Run Configuration in PyCharm.

For the conda environment.

Install Anaconda.

Open an Anaconda cmd.

We will now need to create the environment using these instructions: <u>https://docs.anaconda.com/anacondaorg/user-guide/tasks/work-with-environments/</u>

We have uploaded our environment into our personal account, (sooodos/quantum) and we can activate it with the following command.

conda create --name my_env_name sooodos/quantum

After the installation finishes, go back to PyCharm. In the bottom-right corner click on the Project Interpreter box.

Choose the "Add Interpreter..." option.

From the pop-up window choose "Conda Environment" and then click on the "Existing environment" option.

Now you need to search for the python.exe script responsible for the Interpreter.

By default, it should be located in C:\Users\NameOfUser\Anaconda3\envs\ my_env_name \python.exe.

Save changes.

Now we need to add a Run Configuration from PyCharm.

Go to the upper-right corner of the window, the Run Configurations box is located to the left of the "run" icon.

Click on it and then click "Edit Configurations".

On the pop-up window click the "+" icon and choose "Python". Now we need to add the script path of manager.py in order to create our Run Configuration. Just click on the folder icon which is on the far-right corner of the "Script path" place holder and find manager.py from the project hierarchy.

Click Apply and OK.

Now add the token generated by your IBM Account and insert into the account_details.py file in the corresponding token placeholder.

QCLG is now operational.